

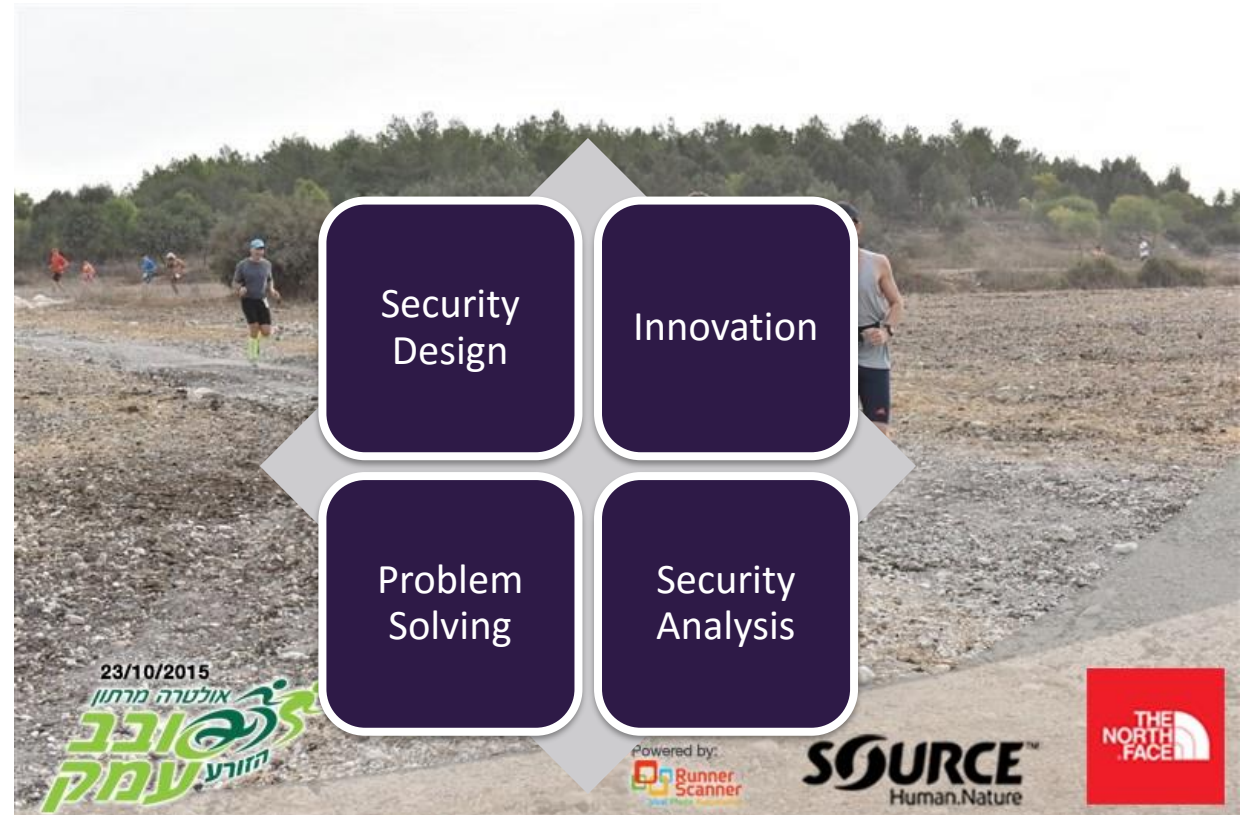
Bar-Mitzva Attack

Breaking SSL with 13-Year Old RC4 Weakness

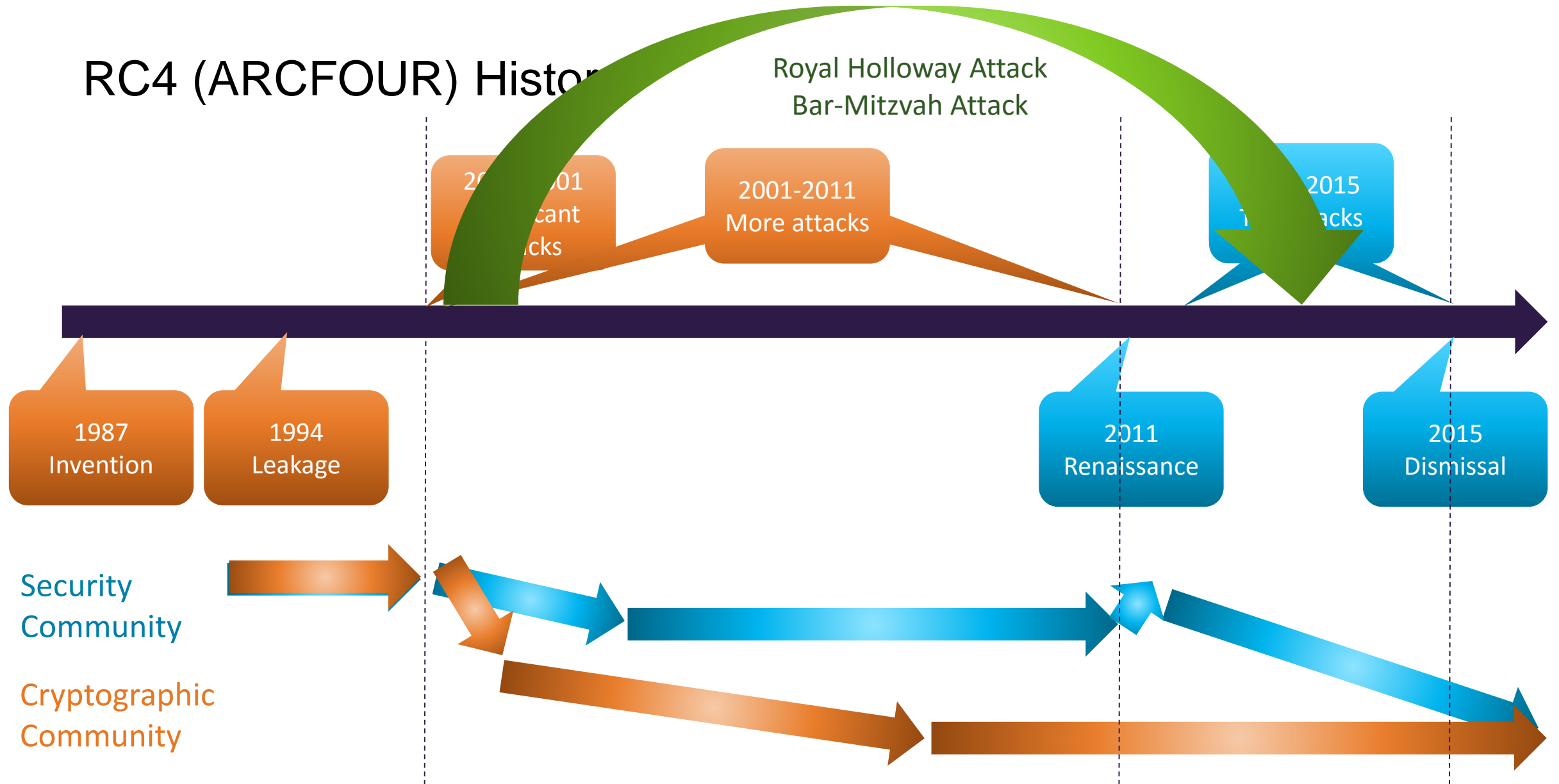
Itsik Mantin

About Myself

- Ultra-Marathoner
- Director of Security Research at Imperva
- Application Defense Center (ADC)
- 16 years in various security domains
 - DRM systems, Web applications, Automotive systems, Insider threats, Cryptography and Cryptanalysis
- M. Sc. in Applied Math and Computer Science from the Weizmann institute
 - Cryptanalysis research with professor Adi Shamir
- <https://www.linkedin.com/in/imantin>

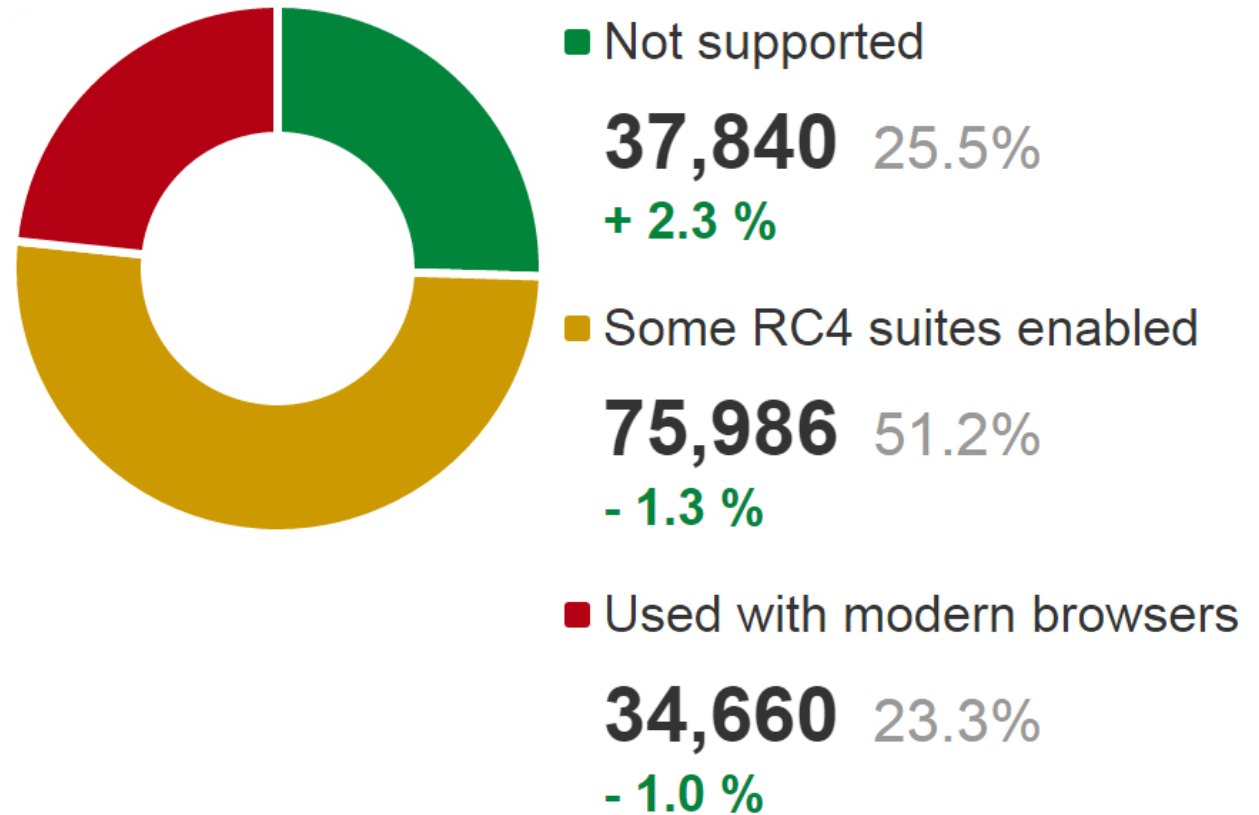


RC4 (ARCFOUR) History



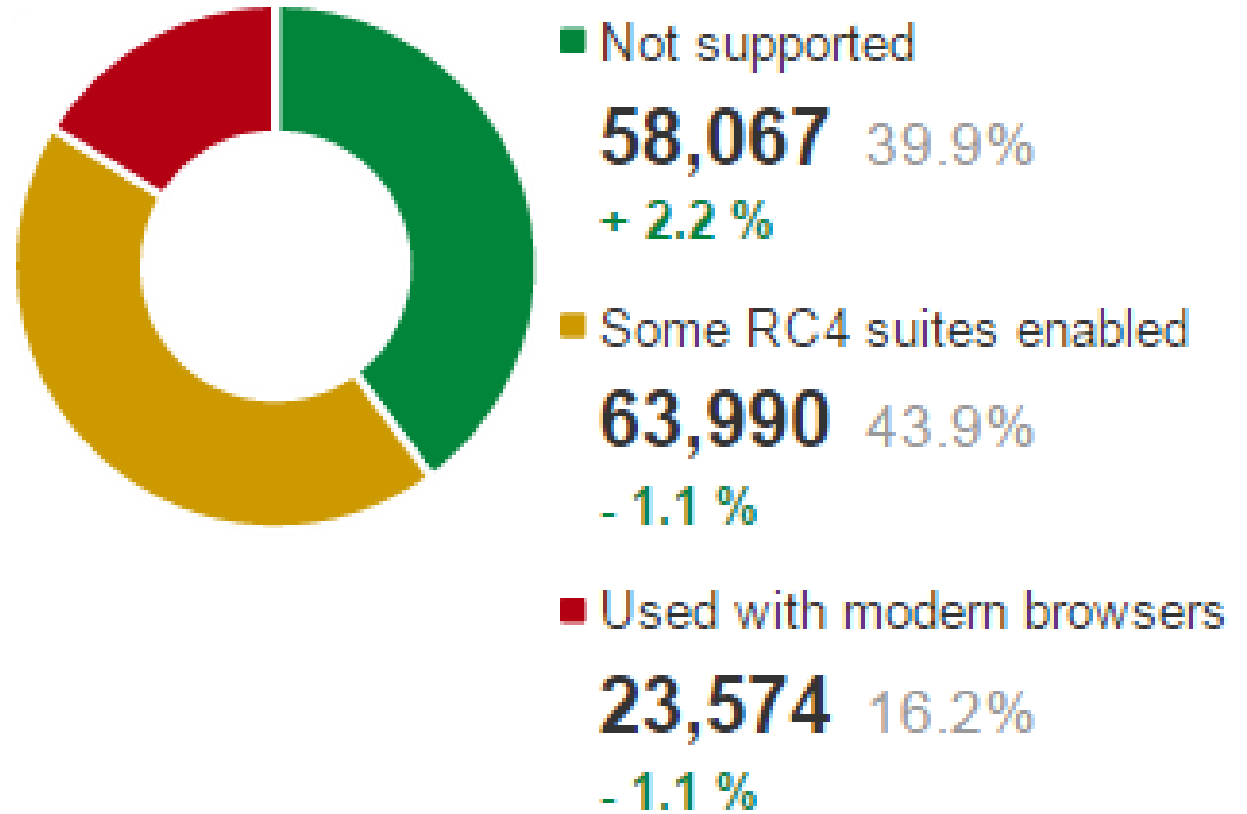
RC4 Usage in TLS

- 150K sites, SSL-Pulse
- March 9, 2015



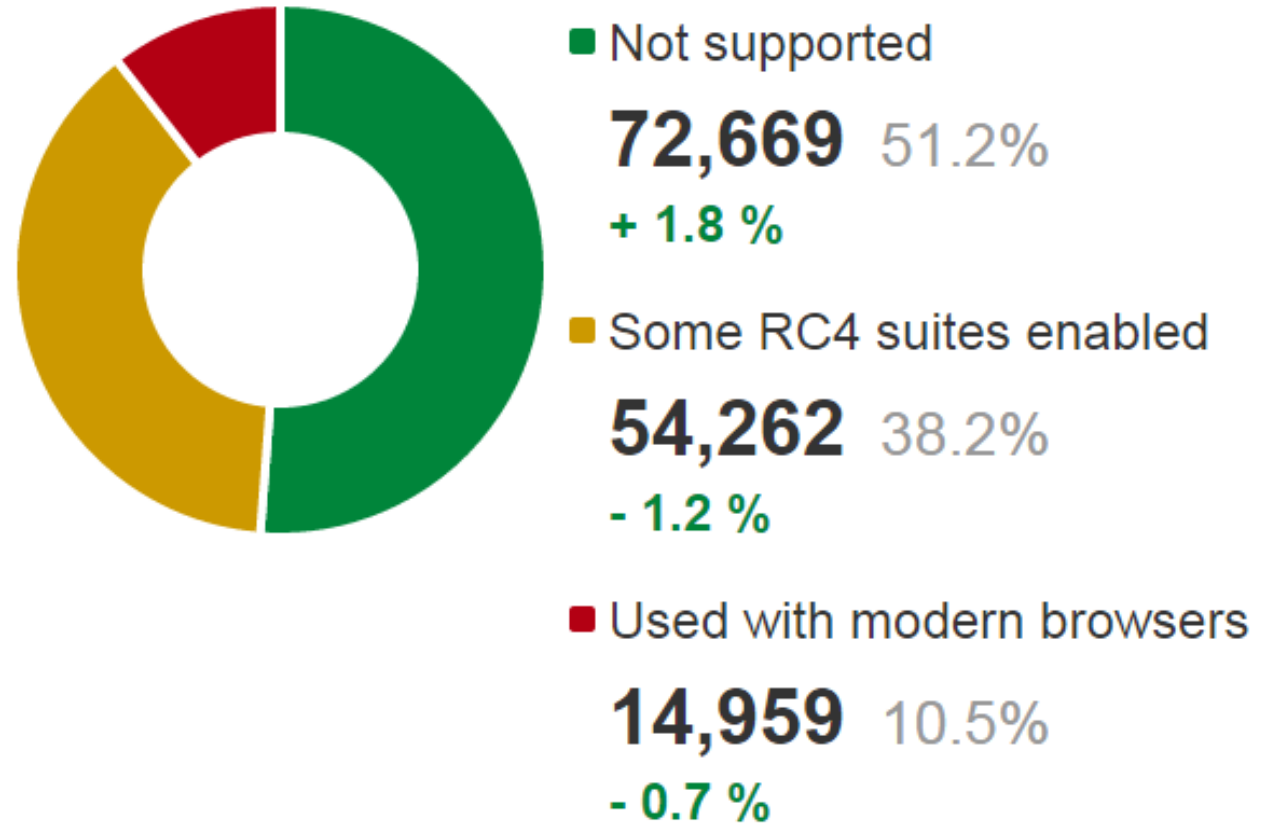
RC4 Usage in TLS

- 150K sites, SSL-Pulse
- July 15, 2015



RC4 Usage in TLS

- 150K sites, SSL-Pulse
- Dec 23, 2015



The Future of RC4

The screenshot shows the InfoWorld website interface. At the top, there is a navigation bar with categories: App Dev, Cloud, Data Center, Mobile, Open Source, Security, Deep Dives, Reviews, and Resources/White Papers. The InfoWorld logo is prominently displayed on the left, with a 'Most Popular:' dropdown menu and social media icons on the right. Below the navigation, the breadcrumb 'Home > Security' is visible. The article is part of the 'INFOWORLD TECH WATCH' series by Fahmida Y. Rashid. The main headline reads 'Google, Mozilla, Microsoft browsers will dump RC4 encryption'. Below the headline is a large image of browser logos (Chrome, Firefox, Edge, Internet Explorer) mounted on metal barrels. To the right of the image is a 'MORE LIKE THIS' section with a link to an article on IDG Answers. At the bottom right, there is a sponsored advertisement for Microsoft Cloud with the text: 'Building Trust & Control in the Cloud. In 2015, business continuity is the top goal of cloud investments, replacing lower TCO as...'

Why Bar Mitzvah?



- The Invariance Weakness

- ***“Weaknesses in the key scheduling algorithm of RC4”***.
Fluhrer, Mantin, and Shamir (Selected Areas of Cryptography, 2001)
- ***“Analysis of the stream cipher RC4”***. Mantin (My M. Sc. Thesis, 2001)

1

On TLS

TLS Objectives

Mutual Authentication

- Usually only Server authentication is used

Data Protection

- Data Integrity
- Data Confidentiality

Passive Attacker (Sniffing)



Man-in-the-Middle Attacker (MitM)



alice.wonder@gmail.com
Alice123!



alice.wonder@gmail.com
Alice123!



TLS Security

Cipher attacks

- BEAST (2011)
- Royal Holloway (2013)

Compression attacks

- CRIME (2012)
- TIME (2013)
- BREACH (2013)

Downgrade attacks

- False Start (2012)
- POODLE (2014)
- FREAK (2014)

Padding Oracle attacks

- Lucky13 (2013)

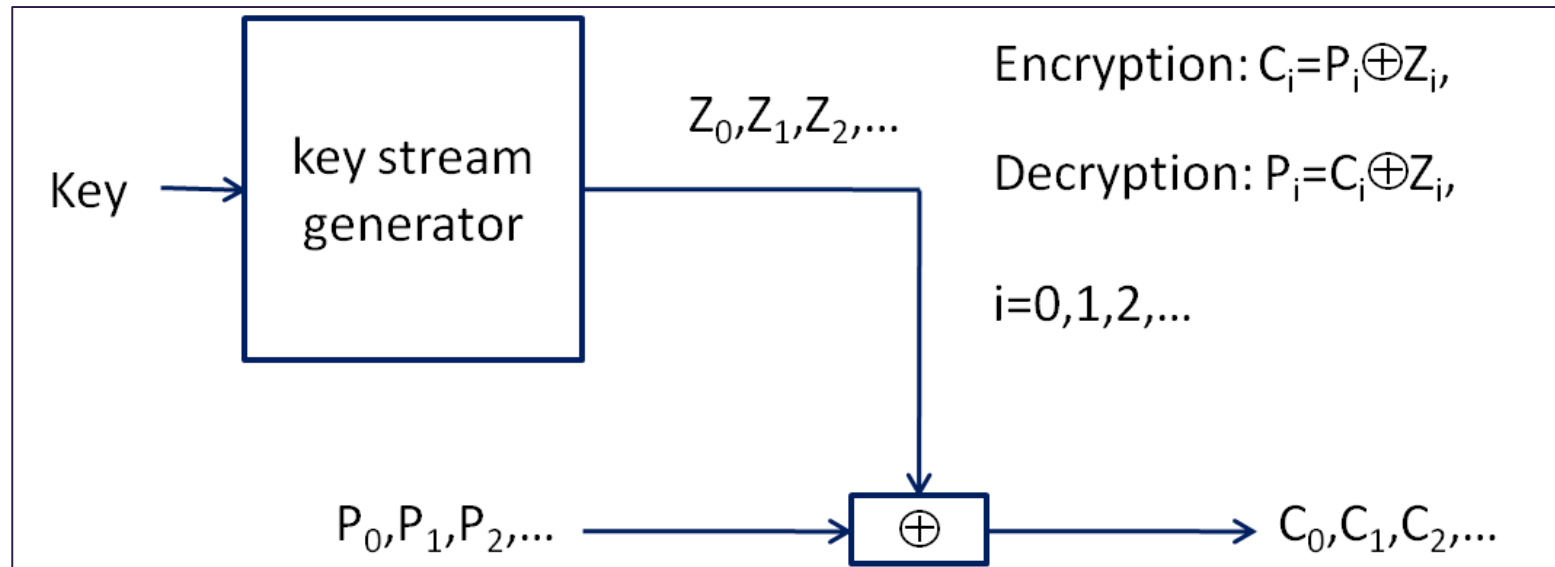
Implementation attacks

- Heartbleed (2014)

2

On RC4

Stream Ciphers



Keystream randomness = plaintext security

RC4

- Rivest Code 4
- The most popular Stream Cipher for more than 25 years
- Details kept secret until the WEP attack in 2001

RC4 Algorithm

Key Scheduling Algorithm (KSA)

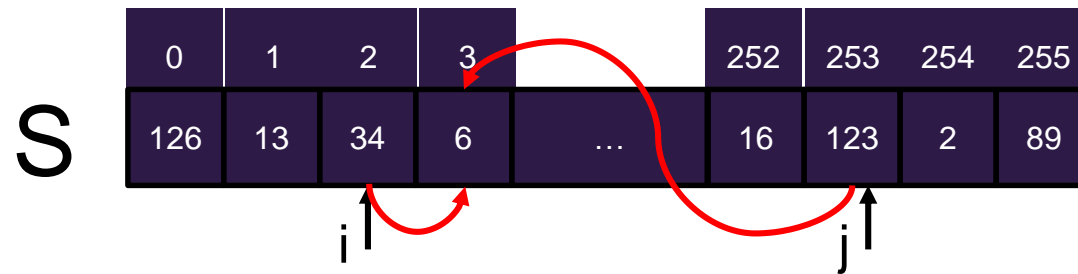
Pseudo-Random Generation Algorithm (PRGA)

All operations are mod 256

```
KSA(K):  
  j = 0  
  S = [0, 1, 2, ..., 255]  
  for i = 0..255  
    j = j + S[i] + K[i mode |K|]  
    S[i] ↔ S[j]
```

```
PRGA(S0):  
  i = 0  
  j = 0  
  S = S0  
  While bytes are needed:  
    i = i + 1  
    j = j + S[i]  
    S[i] ↔ S[j]  
    Emit S[S[i]+S[j]]
```

RC4 Algorithm



KSA(K):

$j = 0$

$S = [0, 1, 2, \dots, 255]$

for $i = 0..255$

$j = j + S[i] + K[i \text{ mode } |K|]$

$S[i] \leftrightarrow S[j]$

PRGA(S_0):

$i = 0$

$j = 0$

$S = S_0$

While bytes are needed:

$i = i + 1$

$j = j + S[i]$

$S[i] \leftrightarrow S[j]$

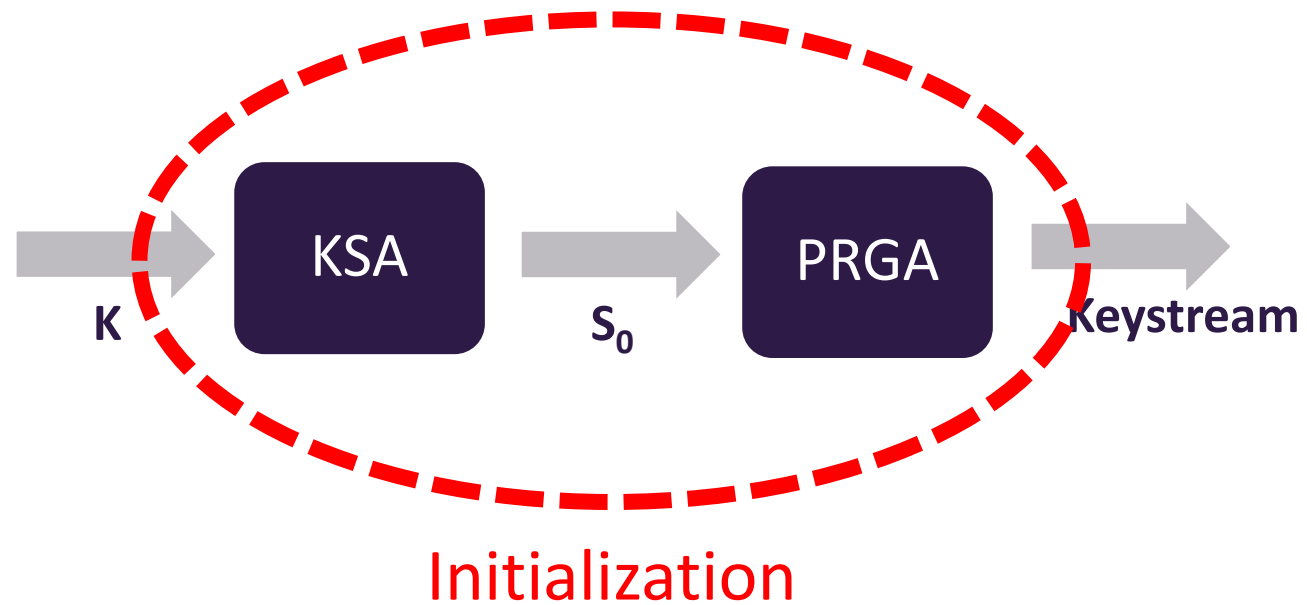
Emit $S[S[i]+S[j]]$

RC4 (In)Randomness

- RC4 in **NOT** pseudo-random
 - 2^{30} distinguishing algorithm
Fluhrer-McGrew, 2000
Patterns used in **2013** to attack TLS (the Royal-Holloway attack)
 - 2^{26} byte distinguishing algorithm
Mantin, 2005
Patterns used in **July 2015** to attack WPA-TKIP and TLS
 - 2^{45} Prediction algorithm
Mantin, 2005

RC4 Initialization

The weakest link of RC4 since 2001



RC4 Initialization

Keystream biases

- The second-byte bias (Mantin-Shamir, 2001)
- Many others

Initial permutation biases

- My thesis 2001, Mironov 2002

Key-keystream correlations

- The IV Weakness and the WEP Attack (Fluhrer-Mantin-Shamir, 2001)
- Enhanced WEP Attack I (Mantin, 2005)
- Enhanced WEP Attack II (Tews-Weinmann-Pyshkin, 2007)
- More Key-keystream correlations (Klein, 2005)
- **The Invariance Weakness (Fluhrer-Mantin-Shamir, 2001)**

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The Invariance Weakness

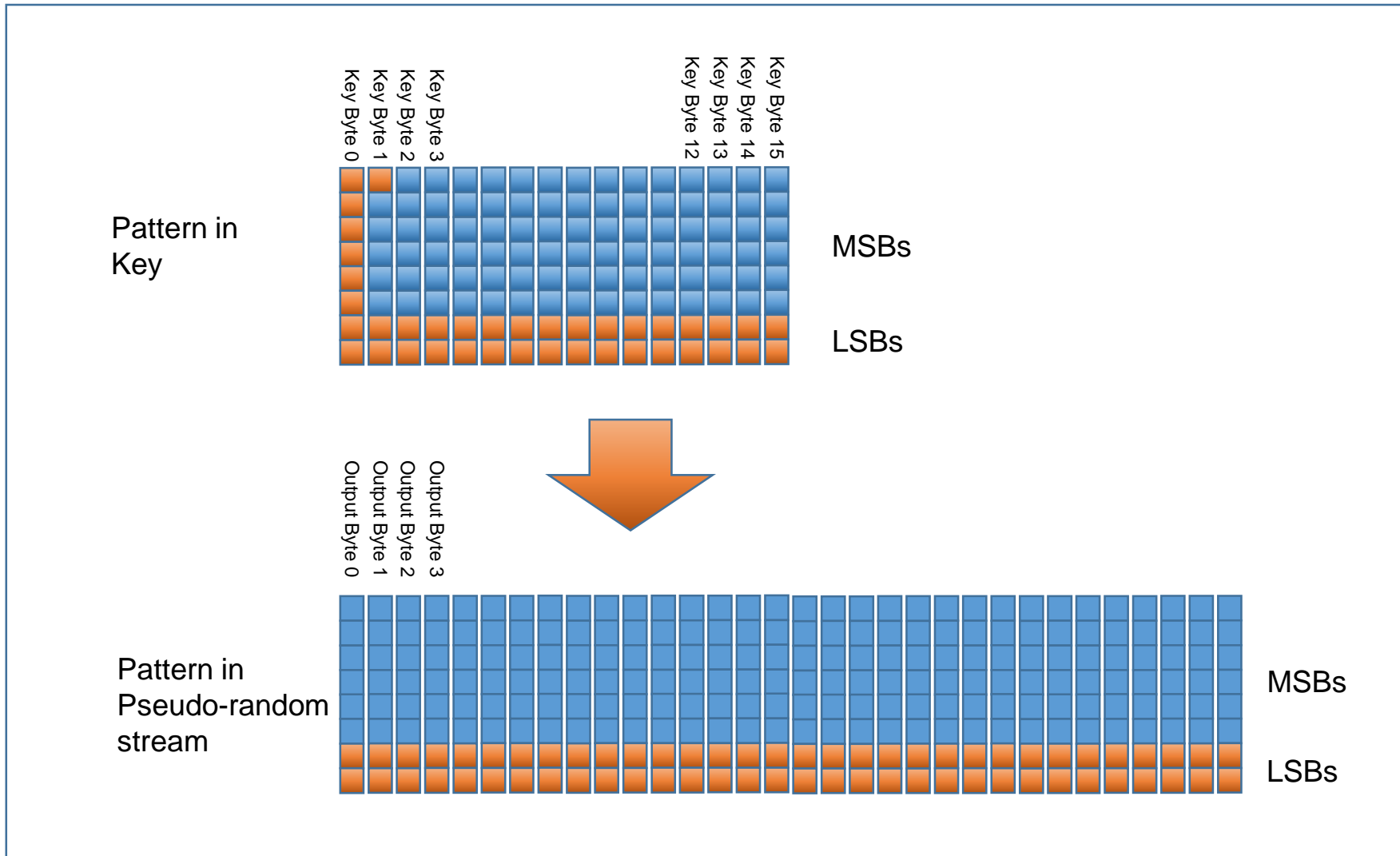
The Invariance Weakness

- The neglected counterpart of the IV Weakness
- Left in the shadows for 13 years

- RC4 weak keys
 - **Huge** class of keys (2^{-24} fraction for 128bit keys)

 - Bad mixing of the key with the permutation.
Permutation parts remain **intact**

Key Patterns



Plaintext Leakage



Weak Key Classes

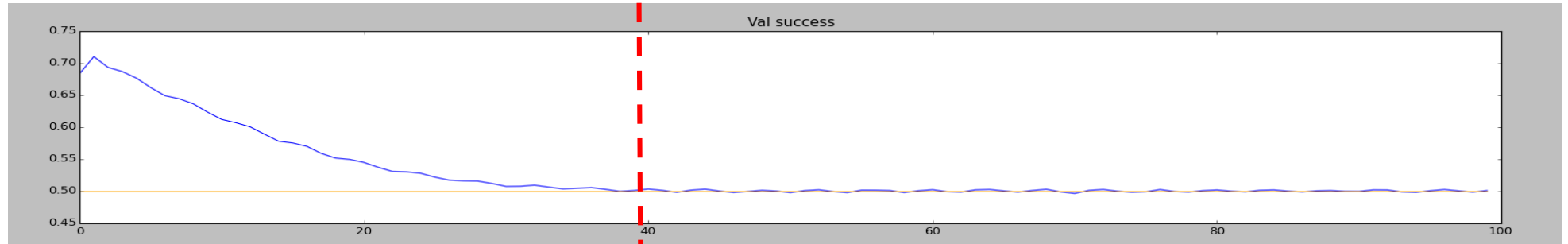
# LSBs	Applicability	Class Fraction (8-byte key)	Class Fraction (16-byte key)
1	Keys with even number of bytes	2^{-16}	2^{-24}
2	Keys with number of bytes that divides 4	2^{-23}	2^{-39}
3	Keys with number of bytes that divides 8	2^{-30}	2^{-54}
4	Keys with number of bytes that divides 16	2^{-37}	2^{-69}

Plaintext Leakage

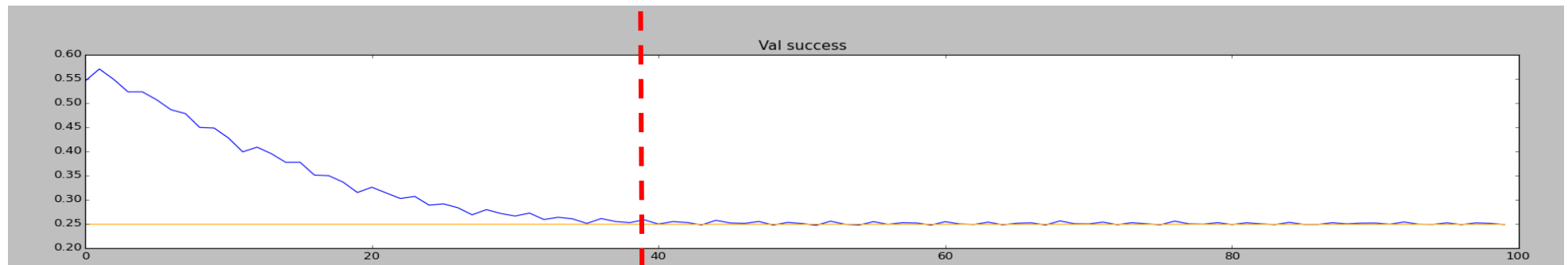
- When a weak key is used, “many” plaintext bit leak
- Q1: Can we tell when that happens?
 - Yes, when plaintext patterns exist
- Q2: How many bits?

Leakage Statistics

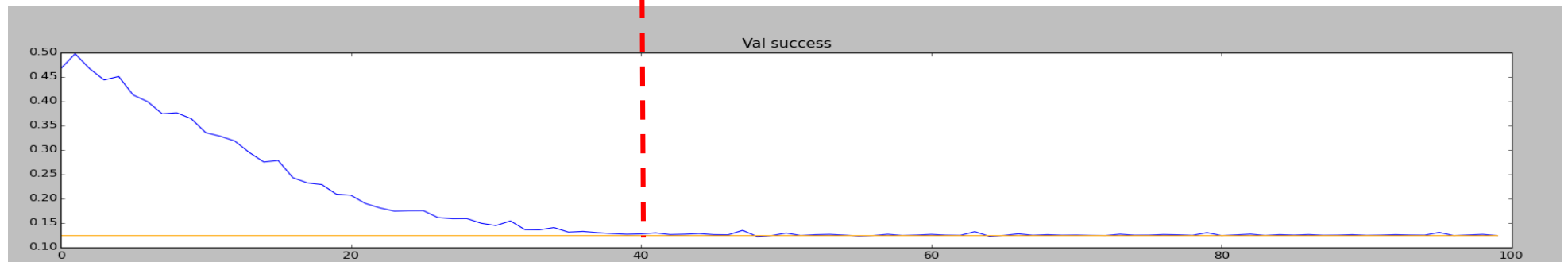
q=1



q=2



q=3

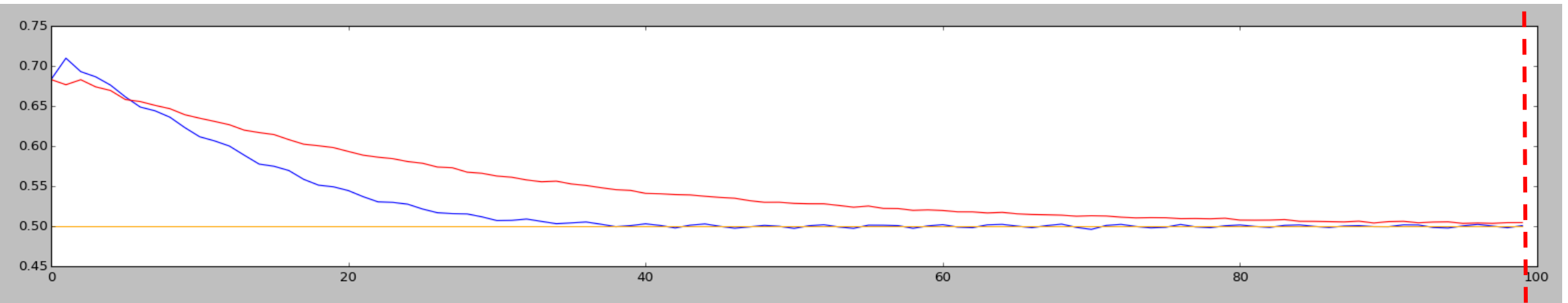


Diff-Based Leakage

- The permutation is ruined with the keystream generation
- Bit prediction gets out of sync when j hits a “ruined” part
- Switch to diff

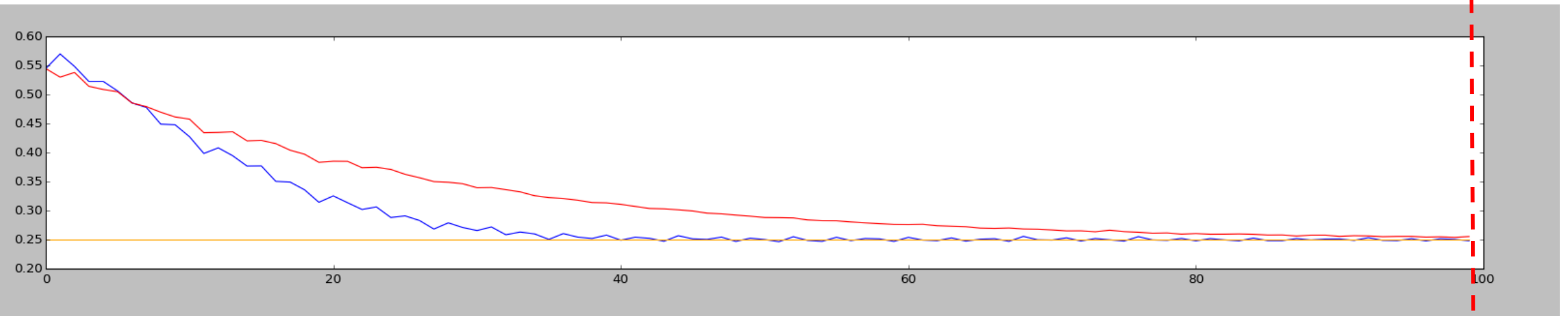
Diff-Based Leakage (q=1)

~100



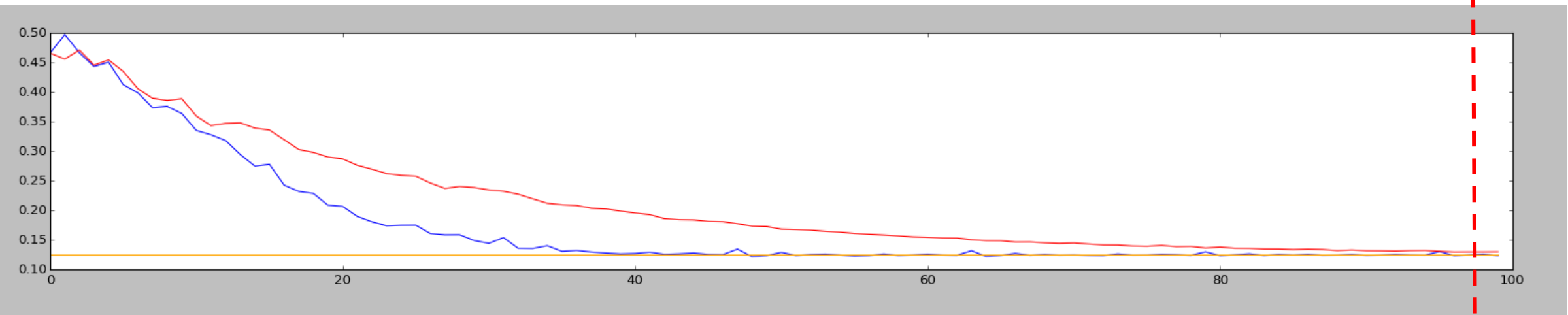
Diff-Based Leakage (q=2)

~100



Diff-Based Leakage (q=3)

~100



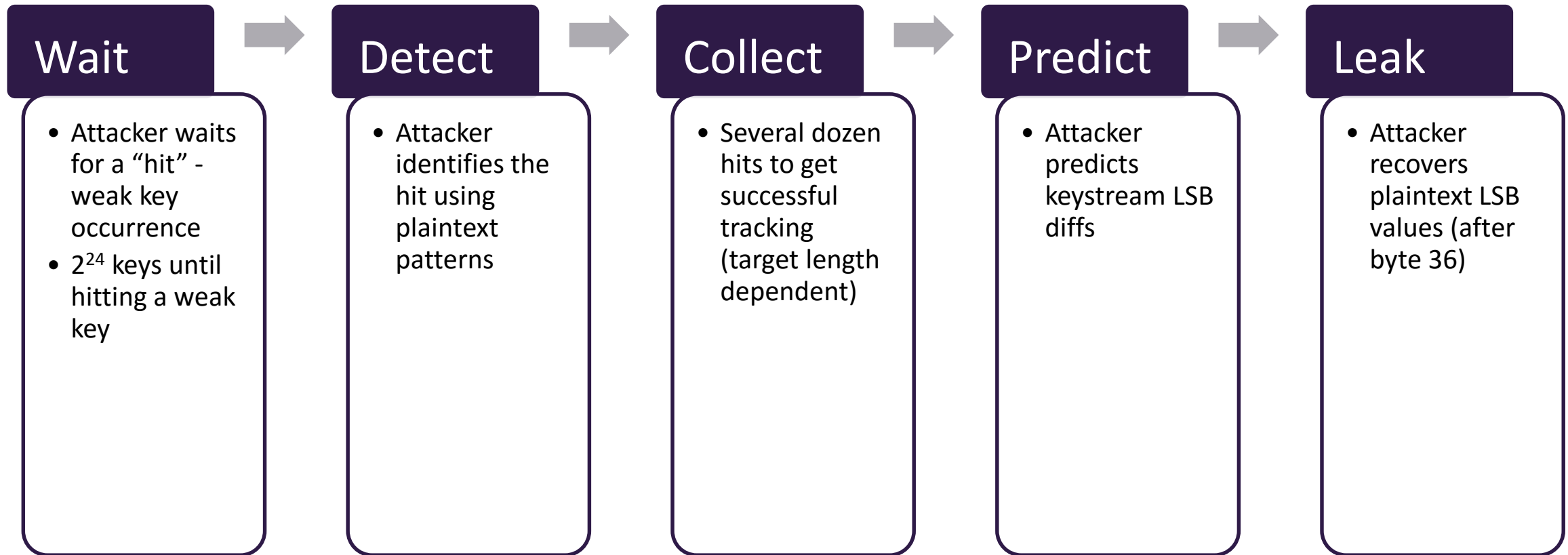
The Leakage

- Using the 1-Class
 - 1st diff LSB is guessed correctly with probability 0.68
 - 37th diff LSB is guessed correctly with probability of 0.546
 - 100th diff LSB is guessed correctly with probability of 0.503
- Pattern tracking is possible for
 - 37 bytes with 1/22 advantage
 - 68 bytes with 1/64 advantage
 - 100 bytes with 1/330 advantage
- **First 100 LSBs are exposed to leakage**

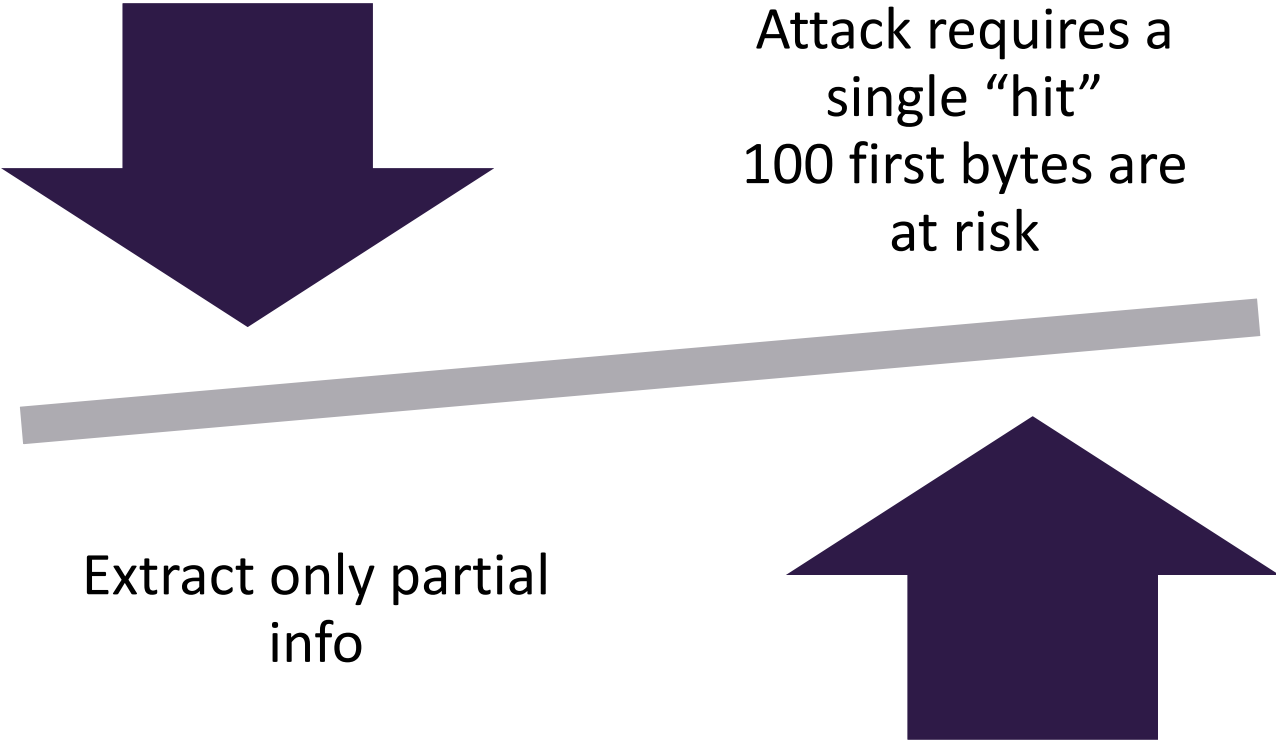
4

The Attacks

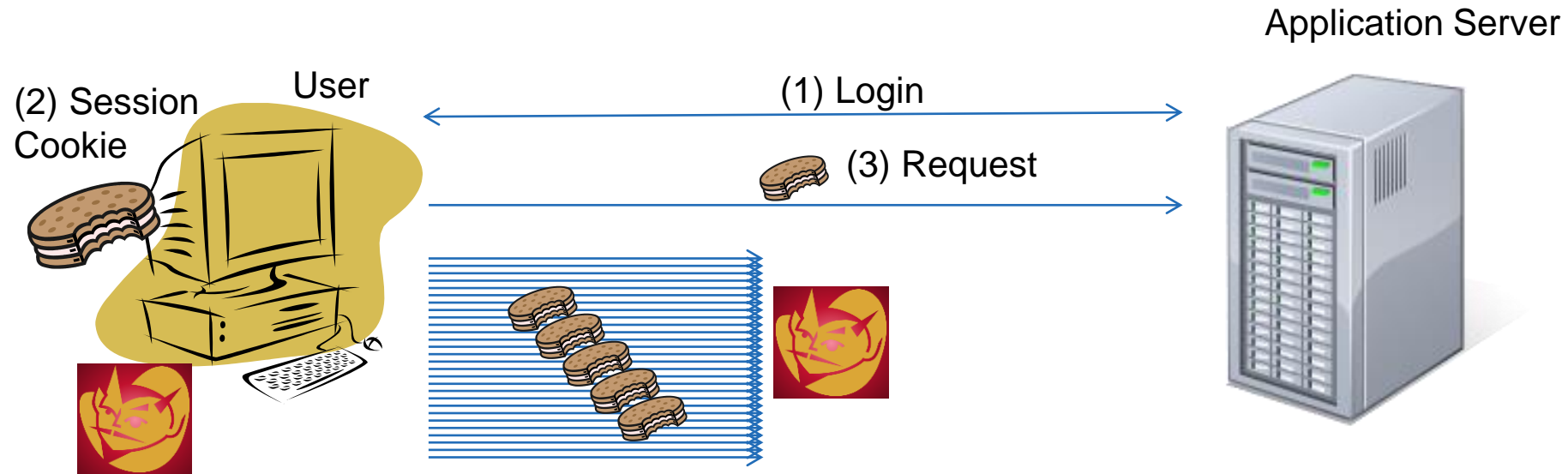
The Attack Basic Scenario



Attack Unique Characteristics



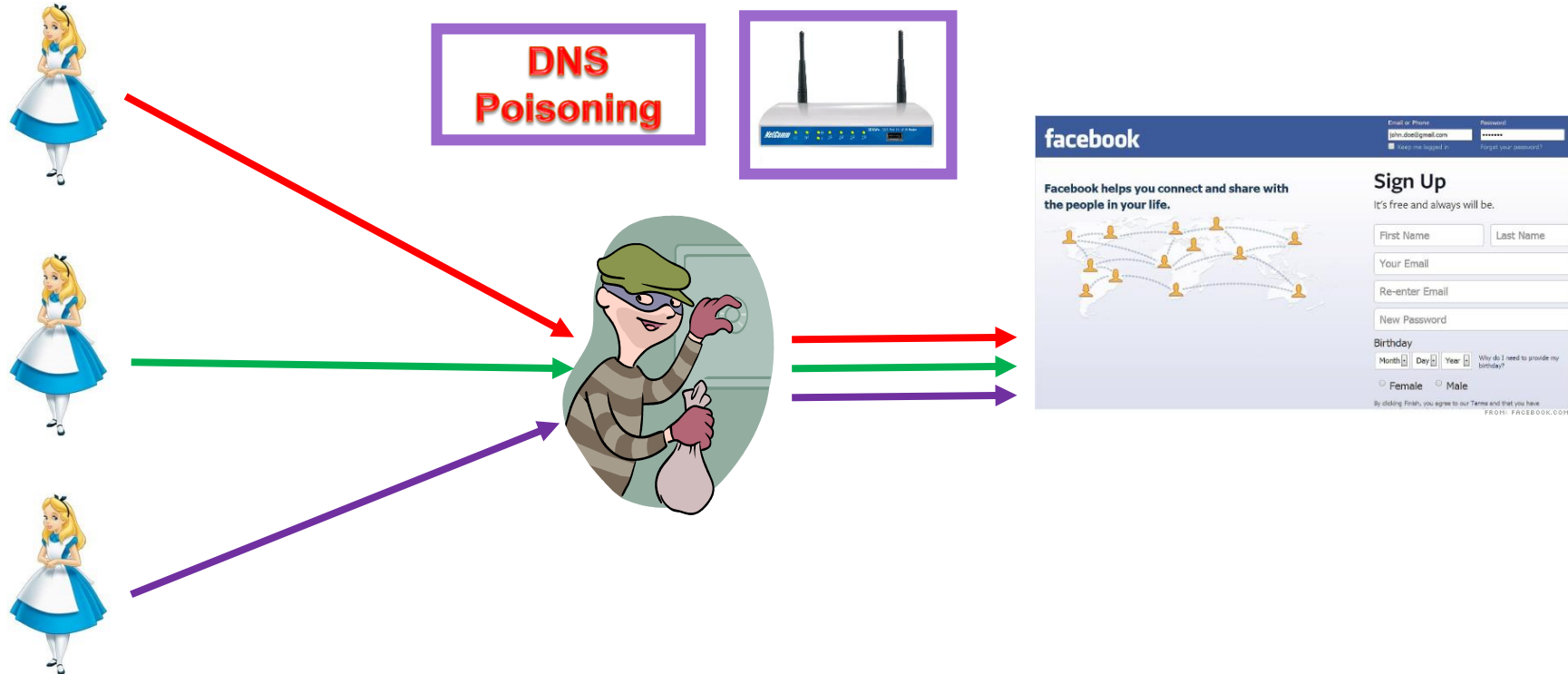
BEAST-like Attack



- 1 billion connections required
- **Insensitive to Resets**

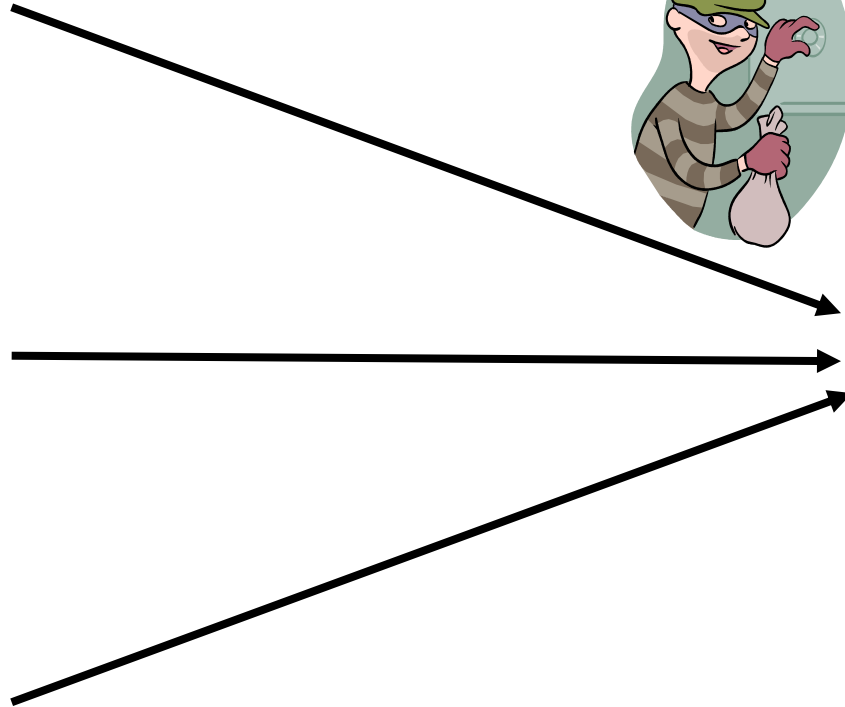
Group Attack

Attack requires a single hit
Pool of Potential Victims



Non-Targeted *Passive* Attack

Attack requires a single “hit”
Pool of Potential Victims



1 Billion Connections?

- Facebook has 890 million DAU (Daily Active Users)
- Most login more than once a day



A Concerning Fact

- Every time you send a secret over TLS/RC4 connection
 - You have a 1:16 million chance to get a bad key
 - You have a 1 in a billion chance to get unlucky and leak a significant portion of your secret
- Small numbers, but definitely not negligible
- RC4 stats (March 2015): 30% of Internet TLS connections

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Conclusion

Summary

- The Invariance Weakness of RC4 can be used to mount new attacks on TLS
- The ***Reset Insensitivity*** nature of the attack opens the door to new attack scenarios
- First passive attack on TLS

Conclusions

- RC4 is not a secure cipher (old news)
- The initialization mechanism of RC4 is very weak (old news)
- The impact of these facts on the (In)Security of systems using RC4 is underestimated (today this is also old news)

More info

- Report at Imperva ADC site:
 - <http://www.imperva.com/DefenseCenter/HackerIntelligenceReports>
- Blackhat materials (white paper and presentation)
 - <https://www.blackhat.com/docs/asia-15/materials/asia-15-Mantin-Bar-Mitzvah-Attack-Breaking-SSL-With-13-Year-Old-RC4-Weakness-wp.pdf>
 - <https://www.blackhat.com/docs/asia-15/materials/asia-15-Mantin-Bar-Mitzvah-Attack-Breaking-SSL-With-13-Year-Old-RC4-Weakness.pdf>
- Wiki of the attack
 - https://en.wikipedia.org/wiki/Bar_mitzvah_attack

Verschlüsselung: Deutlich verbesserte Angriffe auf RC4

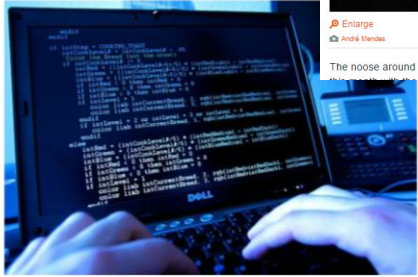
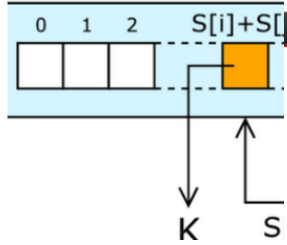
27.03.2015 15:29 Uhr - Jürgen Schmidt

加密等於虛設, SSL、TLS 有存在 13 年的大漏洞

作者: Umire Pro | 發布日期: 2015 年 04 月 06 日 00:00 | 台灣 網路 資訊安全

7 2,404

分享到 Facebook 傳送到 Messenger



SSL 與 TLS 多年來被業界沿用為數據傳輸中的加密標準, 但日前就有資安人員揭露指, 原來上述兩種加密方式依然存在著根本上的漏洞, 被發現的漏洞更出現長達 13 年之久, 用戶敏感資料完全沒保障。

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Noose around Internet's TLS system tightens with 2 new decryption attacks

Exploits pluck passwords and other sensitive data out of encrypted data streams.

by Dan Goodin - Mar 27, 2015 2:22am EET



The noose around the neck of the Internet's most widely used encryption scheme got a little tighter

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ZUM VORTEILSANGEBOT

ZDNet / Sicherheit / Sicherheitsmanagement
Forscher demonstrieren Schwachstelle bei RC4-Verschlüsselung für SSL-Verbindungen
von Kai Schirmer am 27. März 2015, 18:23 Uhr

Der Sicherheitspezialist Imperva hat einen neuen Hacker Intelligence Initiative (HII) Report veröffentlicht. Demnach hat sein Forschungsteam das Application Defense Center (ADC) neue Angriffsschwachstellen in dem verbreiteten Transport Layer Security (TLS/SSL)-Zertifikat entdeckt.

IMPERVA Whitepaper

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ATTACKS/BREACHES

SSL/TLS Suffers 'Bar Mitzvah Attack'

Researcher at Black Hat Asia shows how attackers could abuse a known-weak crypto algorithm to steal credentials and other data from encrypted communications.

SSL/TLS encryption once again is being haunted by an outdated and weak feature long past its prime: a newly discovered attack exploits a weakness in the older, less secure RC4 encryption algorithm option in SSL/TLS that's still supported in many browsers and servers.

Itsk Martin, director of security research with Imperva, at Black Hat Asia in Singapore today will detail how an attacker could sniff credentials and other information during an SSL session in an attack he named the "Bar Mitzvah Attack" after 13-year-old weaknesses in the algorithm it abuses. The attack is a nifty reminder that the RC4 algorithm, long known to be breakable, should

EDUCATIONAL RESOURCES

- Beyond Malware: Detecting the Undetectable
- Indicators of Attack vs. Indicators of Compromise
- Next Generation Endpoint Protection Case Study: Cardinal Innovations Healthcare
- Cyber Intrusion Services Report
- Crowdstrike Services Brochure

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