

## Lecture 9

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Last lecture:

- An alternative proof of Schützenberger's Theorem
- Detailed discussion of  $DU$ -index.

## Congruences for density functions

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- **The Little Fermat Theorem:** A combinatorial proof
- **Gessel's Theorem:** Bounded degree Gessel classes of (directed) graphs
- Classes of relational structures of bounded degree and finite  $DU$ -index (Fischer-M., 2003).
- **The Specker Blatter Theorem:** Classes of relational structures with relations of arity 2 and finite Specker index.

## Fermat's Theorem: $a^p = a \pmod p$

### A combinatorial proof

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- Look at the wheel  $W_p$  with  $p$  a prime. The edges connecting the center to the circle are called **spikes**.
- Let  $f : Spikes \rightarrow [a]$  a coloring of the spikes with  $a$  colors. There are  $a^p$  such colorings.
- Two colorings  $f_1$  and  $f_2$  are equivalent if the wheel can be turned from an  $f_1$  coloring to an  $f_2$  coloring.
- All monochromatic colorings form singleton equivalence classes. There are  $a$  such classes.
- All other equivalence classes have  $p$  members.
- Hence  $a^p - a \equiv 0 \pmod p$ .

Q.E.D.

## Gessel's Theorem (1984)

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### Theorem 1 (I. Gessel 1984)

If  $\mathcal{C}$  is a Gessel class of directed graphs of degree at most  $d$  with density function  $d_{\mathcal{C}}(n)$ , then

$$d_{\mathcal{C}}(m + n) \equiv d_{\mathcal{C}}(m) \cdot d_{\mathcal{C}}(n) \pmod{\frac{m}{\ell}}$$

where  $\ell$  is the least common multiple of all divisors of  $m$  not greater than  $d$ .

## Gessel's Theorem (reformulated as linear modular recurrence)

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### Theorem 2 (I. Gessel 1984)

If  $\mathcal{C}$  is a Gessel class of (directed) graphs of degree at most  $d$ , then for every  $m$  there is a  $q = q(d, m) \in \mathbb{N}$  such that

$$f_{\mathcal{C}}(n + q) \equiv \sum_{i=0}^{q-1} b_i \cdot f_{\mathcal{C}}(n + i) \pmod{m}$$

with  $b_0 = f_{\mathcal{C}}(q)$  and  $b_i = 0 : 1 \leq i \leq q - 1$ .

Furthemore,  $q = d!m$ .

*Exercise:* Can you make  $q$  smaller?

## Orbit of a structure

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Let  $\mathcal{G} \subseteq \mathcal{S}_n$  be a subgroup of the full permutation group of  $[n] = \{0, \dots, n-1\}$ .

For a  $\tau$ -structure  $\mathfrak{A}$  over the universe  $A = [n]$

- the orbit  $Orb_{\mathcal{G}}(\mathfrak{A})$  is the set of different labeled structures  $\sigma(\mathfrak{A})$  obtained using relabelings from  $\mathcal{G}$ , i.e.

$$Orb_{\mathcal{G}}(\mathfrak{A}) = \{\sigma(\mathfrak{A}) : \sigma \in \mathcal{G}\}$$

- If  $\mathcal{G} = \mathcal{S}_n$ , we omit it and write  $Orb(\mathfrak{A})$
- $Aut_{\mathcal{G}}(\mathfrak{A})$  is the set of  $\tau$ -automorphisms of  $\mathfrak{A}$  which are in  $\mathcal{G}$ .

As  $Aut_{\mathcal{G}}(\mathfrak{A}) \subseteq \mathcal{G}$  is a subgroup, we have the fundamental identity

### Proposition 3

For a  $\tau$ -structure  $\mathfrak{A}$  over the universe  $A = [n]$  and  $\mathcal{G}$  a subgroup of  $\mathcal{S}_n$  we have

$$|Aut_{\mathcal{G}}(\mathfrak{A})| \cdot |Orb_{\mathcal{G}}(\mathfrak{A})| = |\mathcal{G}|$$

## Gessel's Lemma

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Here the group is the cyclic group  $\mathbb{Z}_m$  of cyclic permutations of  $[m]$ .

$[m]$  is the set  $\{1, \dots, m\}$ .  $m + [n]$  is the set  $\{m + 1, \dots, m + n\}$ .

Let  $G = (V(G), E(G))$  be a graph and  $A, B \subseteq V(G)$ .

$A$  is **adjacent to**  $B$  if there is an edge between  $A$  and  $B$ .

### Lemma 4

*Let  $G$  a graph on  $[m + n]$  with degree at most  $d$ .*

*Let  $\ell$  be the least common multiple of all divisors of  $m$  not greater than  $d$ .*

*Assume in  $G$  the sets  $[m]$  and  $m + [n]$  are adjacent.*

*Then  $|\text{Orb}_{\mathbb{Z}_m}(G)|$  is a multiple of  $\frac{m}{\ell}$ .*

Proof: Suppose  $(i, j) \in E(G)$  with  $i \in [m]$  and  $j \in m + [n]$ .

Let  $g \in \text{Aut}_{\mathbb{Z}_m}(G)$ . Then  $(g(i), g(j)) = (g(i), j) \in E(G)$ .

So  $|\text{Aut}_{\mathbb{Z}_m}(G)| \leq d$ , and  $|\text{Aut}_{\mathbb{Z}_m}(G)|$  divides  $\ell$ .

Now the lemma follows from Proposition 3.

Q.E.D.

## Proof of Gessel's Theorem (Theorem 1)

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**Step 1:** As Gessel classes are closed under components,  $d_{\mathcal{C}}(m) \cdot d_{\mathcal{C}}(n)$  counts the number of graphs  $G$  on  $[m+n]$  where there is no edge between  $[m]$  and  $m+[n]$ .

**Step 2:**  $d_{\mathcal{C}}(m+n) - d_{\mathcal{C}}(m) \cdot d_{\mathcal{C}}(n)$  counts the number of graphs for which Gessel's Lemma applies.

Therefore  $d_{\mathcal{C}}(m+n) - d_{\mathcal{C}}(m) \cdot d_{\mathcal{C}}(n) \equiv 0 \pmod{\frac{m}{\ell}}$

Q.E.D.



## Bounded $DU$ -index and bounded degree

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### **Theorem 5 (E. Fischer and J.A. Makowsky, 2002)**

*Let  $\mathcal{P}$  be a property of  $\tau$ -structures, with finite  $DU$ -index and all its members of bounded degree  $d$ . Then*

- *$d_{\mathcal{P}}(n)$  satisfies a modular recurrence relation for every  $m$ .*
- *Furthermore, if additionally all the models in  $\mathcal{P}$  are connected, (hence the  $DU$ -index is 2, the function  $f_{\mathcal{P}}$  satisfies the trivial recurrence relations for every  $m$ .*

Example  $\overline{EQ_2CLIQUE}(A)$  shows that bounded degree cannot be dropped, even for  $DU$ -index 2 (and connected structures).

For structures of unbounded degree one needs a stronger assumption, the finiteness of the Specker-index, to be discussed later (if time permits).

## Density functions and orbits

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Let  $\mathcal{P}$  be a class of  $\tau$ -structures and  $\mathcal{P}^n$  the class of structures with universe  $[n]$  in  $\mathcal{P}$ .

We denote by  $Iso(\mathcal{P})$  the  $\tau$ -isomorphisms classes of  $\mathcal{P}$  and denote its equivalence classes by  $[\mathfrak{A}]_{\mathcal{P}}$ .

The orbits  $Orb(\mathfrak{A})$  and the density function  $d_{\mathcal{P}}(n)$  are related by the following formula:

$$d_{\mathcal{P}}(n) = \sum_{[\mathfrak{A}]_{\mathcal{P}} \in Iso(\mathcal{P}^n)} Orb(\mathfrak{A})$$

To show that  $d_{\mathcal{P}}(n) = 0 \pmod{m}$  it suffices to show that for each  $\mathfrak{A} \in \mathcal{P}^n$  we have  $Orb(\mathfrak{A}) = 0 \pmod{m}$ .

## Density functions and orbits via a subgroup

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Let  $\mathcal{G}$  be subgroup of  $\mathcal{S}_n$ .

$\mathcal{G}$  induces an equivalence relation on  $\mathcal{P}$ :

$$\mathfrak{A} \sim_{\mathcal{G}} \mathfrak{A}' \text{ iff there is } \sigma \in \mathcal{G} \text{ with } \sigma(\mathfrak{A}) = \mathfrak{A}'$$

We denote by  $\mathcal{P}/\mathcal{G}$  the set of these equivalence classes in  $\mathcal{P}^n$  and denote its equivalence classes by  $[\mathfrak{A}]_{\mathcal{G}}$ .

### Lemma 6

*The orbits  $Orb_{\mathcal{G}}(\mathfrak{A})$  and the density function  $d_{\mathcal{P}}(n)$  are related by the following formula:*

$$d_{\mathcal{P}}(n) = \sum_{[\mathfrak{A}]_{\mathcal{G}} \in \mathcal{P}^n/\mathcal{G}} Orb_{\mathcal{G}}(\mathfrak{A})$$

To show that  $d_{\mathcal{P}}(n) = 0 \pmod{m}$  it suffices to show that for each  $\mathfrak{A} \in \mathcal{P}^n$  we have  $Orb_{\mathcal{G}}(\mathfrak{A}) = 0 \pmod{m}$ .

## Degrees and orbits

### Lemma 7

Let  $A = [n]$  and  $B \subseteq A$  and  $a \in A - B$ .

Let  $N_a^B$  be the set of neighbors of  $a$  which are in  $B$ .

- Let  $S_B \subseteq S_n$  be the subgroup of permutations  $\sigma$  such that  $\sigma(a) = a$  for every  $a \in A - B$ .
- Let  $G_B^{N_a} \subseteq S_B$  be the subgroup of  $S_B$  which maps  $N_a^B$  onto itself.

Then  $|Orb_{S_B}(\mathfrak{A})|$  is divisible by  $\binom{|B|}{|N_a^B|}$ .

### Proof:

$$|G_B^{N_a}| \cdot \binom{|B|}{|N_a^B|} = |S_B| = |Aut_{S_B}(\mathfrak{A})| \cdot |Orb_{S_B}(\mathfrak{A})| =$$

$$|Orb_{G_B^{N_a}}(\mathfrak{A})| \cdot |Aut_{G_B^{N_a}}(\mathfrak{A})| \cdot \binom{|B|}{|N_a^B|}$$

But we have  $Aut_{G_B^{N_a}}(\mathfrak{A}) = Aut_{S_B}(\mathfrak{A})$ , hence the result.

Q.E.D.

## Choosing special values

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We fix  $m$ , the modulus, and  $d$ , the degree.

We put  $c = d! \cdot m$ .

### Lemma 8

For every  $t \in \mathbb{N}$  and  $0 < d_1 \leq d$  we have that  $m$  divides  $\binom{t \cdot c}{d_1}$ .

**Proof:** Write out the definitions.

$$\begin{aligned} \binom{t \cdot c}{d_1} &= \binom{t \cdot d! \cdot m}{d_1} = \frac{(t \cdot d! \cdot m)!}{(t \cdot d! \cdot m - d_1)! \cdot d_1!} = \\ &= \frac{t \cdot m \cdot d! \cdot \prod_{i=1}^{d_1-1} (t \cdot m \cdot d! - i)}{d_1!} = \\ &= \frac{t \cdot m \cdot d!}{d_1 \cdot (d_1 - 1)!} \cdot \prod_{i=1}^{d_1-1} (t \cdot m \cdot d! - i) = \frac{t \cdot m \cdot d!}{d_1} \cdot \prod_{i=1}^{d_1-1} \frac{(t \cdot m \cdot d! - i)}{i} \end{aligned}$$

Q.E.D

## Connected structures, I

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Recall:  $m$  is the modulus, and  $d$  the degree.  $c = d! \cdot m$ .

### Lemma 9

*For every  $t \in \mathbb{N}$  and every connected  $\mathfrak{A}$  with  $|A| \geq t \cdot c + 1$ , and for every  $B \subseteq A$  with  $|B| = t \cdot c$ , we have that  $|Orb_{S_B}(\mathfrak{A})|$  is divisible by  $m$ .*

**Proof:** For  $t, A, B$  as required, there is  $a \in A - B$  with at  $d_a$  neighbors in  $B$ , and  $1 \leq d_a \leq d$ .

By Lemma 7  $|Orb_{S_B}(\mathfrak{A})|$  is divisible by  $\binom{|B|}{d_a}$ .

Using Lemma 8 with  $|B| = t \cdot c$  and  $1 \leq d_a \leq d$ , we get, it is also divisible by  $m$ .

## Connected structures, II

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Now we use Lemma 6 with  $\mathcal{G} = S_B$  and  $|B| = t \cdot c$ .

But this means that  $B$  has to be fixed **independently** of the particular structure  $\mathfrak{A}$  in  $\mathcal{P}^n$ .

However, as all the  $\mathfrak{A} \in \mathcal{P}^n$  are connected, there is always an  $a \in A - B$  which has neighbors in  $B$ .

Hence, by Lemma 9, for every  $\mathfrak{A} \in \mathcal{P}^n$   
 $|Orb_{S_B}(\mathfrak{A})|$  is divisible by  $m$ .

Q.E.D.

**Remark:** *We did not use a particular  $\mathcal{P}$  of finite DU-index. We only used connectedness (which implies that the DU-index is 2).*

## Disconnected structures, I

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Let  $\mathcal{P}$  be class of structures of bounded degree  $d$  and finite  $D$ -index  $\alpha$ .

Let  $\mathcal{D}_i, i \leq \alpha$  be the  $DU$ -equivalence classes. with respect to  $\mathcal{P}$ .

- All the structures of degree bigger than  $d$  are in one class, say  $\mathcal{D}_0$ .
- $\mathcal{P}$  is also one of the classes.  
(Allowing  $\mathfrak{B}$  to be empty)

$m$  and  $d$  are still fixed and  $c = m \cdot d!$ .

We look now at structures with universe  $[n]$ .

Let  $t \in \mathbb{N}$  and  $B = [t \cdot c]$ .

- Let  $\mathcal{D}_i^0$  be those structures in  $\mathcal{D}_i$  for which there is  $a \in A - B$  which is connected to some  $b \in B$ , and
- $\mathcal{D}_i^1$  be those structures in which no such  $a$  exists.



## Disconnected structures, II

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We want to compute  $d_i(n) = d_{\mathcal{D}_i}$  and  $d_i^j(n) = d_{\mathcal{D}_i^j}$  modulo  $m$  for each  $i \leq \alpha$  and  $j = 0, 1$ .

Clearly,  $d_i(n) = d_i^0(n) + d_i^1(n)$ .

For  $\mathfrak{A} \in \mathcal{D}_i^0$  we apply Lemma 7. Let  $d_a$  be the number of neighbors of  $a$  in  $B$ .

Hence,  $Orb_{S_B}(\mathfrak{A})$  is divisible by  $\binom{t \cdot c}{d_a}$ .

By Lemma 8  $Orb_{S_B}(\mathfrak{A})$  is divisible by  $m$ , hence  $Orb(\mathfrak{A})$  and  $d_i^0(n)$  are divisible by  $m$ .

**Conclusion:**

$$d_i(n) = d_i^0(n) + d_i^1(n) = d_i^1(n) \pmod{m}$$

## Disconnected structures, III

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For  $\mathfrak{A} \in \mathcal{D}_i^1$  we note that  $\mathfrak{A}$  can be uniquely written as

$$\mathfrak{A}_1 \sqcup \mathfrak{A}_2$$

with universes  $A_1 = [t \cdot c]$  and  $A_2 = \{t \cdot c + 1, \dots, n\}$ .

**Fact:** The equivalence class  $[\mathfrak{A}_1 \sqcup \mathfrak{A}_2]$  is uniquely determined by the equivalence classes  $[\mathfrak{A}_1]$  and  $[\mathfrak{A}_2]$ .

Now we put  $t(n) = \lfloor \frac{n-1}{c} \rfloor$  and  $\hat{n} = n \pmod{c}$ .

**Conclusion:** We get, summing over all possibilities

$$d_i(n) = \sum_{j=1}^{\alpha} \mu_{i,j,m,\hat{n}} d_j(t(n)) \pmod{m}$$

## Disconnected structures, IV

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This gives us  $(\alpha \times \alpha)$ -matrices

$$M(m, \hat{n}) = (\mu_{i,j,m,\hat{n}})_{i,j}$$

Let  $\bar{d}(n) = (d_1(n), \dots, d_\alpha(n))^{tr}$ .

For each  $\hat{n} \in \mathbb{Z}_c$  we get now the relationship

$$\bar{d}(n) = M(m, \hat{n}) \cdot \bar{d}(t(n)) \pmod{m}$$

Using the characteristic polynomials  $p_{m,\hat{n}}(\lambda)$  of all the matrices  $M(m, \hat{n})$ , we can now compute the required linear recurrence modulo  $m$ . Q.E.D.

## Graphs of unbounded degree, I

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The case of graphs of unbounded degree has several complications.

- Finite  $DU$ -index of  $\mathcal{P}$  does not suffice.  
We have to make a stronger assumption: **Finite Specker index**.
- The restriction to relations of arity at most 2 is essential.
- The proof, although similar, is considerably more complicated.

## Specker index, survey

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- To define the Specker index of a graph property  $\mathcal{P}$  one defines a binary operation  $subst(H, a, G)$  where in the pointed graph  $H = (V_H, E_H, a)$  the vertex  $a$  is substituted by  $G = (V_G, E_G)$ .

- $DU(\mathcal{P})$ -equivalence is now replaced by  $subst(\mathcal{P})$ -equivalence:  
 $G_1 \sim_{\mathcal{P}} G_2$  iff for all  $H$  and  $a \in V_H$  we have

$$subst(H, a, G_1) \in \mathcal{P} \text{ iff } subst(H, a, G_2) \in \mathcal{P}$$

- The Specker index of  $\mathcal{P}$  is the number of  $subst(\mathcal{P})$ -equivalence classes.

### Observation:

- The  $DU$ -index of  $\mathcal{P}$  is always smaller or equal The Specker-index of  $\mathcal{P}$ .
- $MSOL$ -definability of  $\mathcal{P}$  implies finite Specker index.

## Pointed structures

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A pointed  $\bar{R}$ -structure is a pair

$$(\mathfrak{A}, a)$$

with

- $\mathfrak{A}$  an  $\bar{R}$ -structure and
- $a$  an element of the universe  $A$  of  $\mathfrak{A}$  or
- $a$  an element not in  $A$ .

In  $(\mathfrak{A}, a)$ , we speak of the structure  $\mathfrak{A}$  and the *context*  $a$ .

## Substitution of structures, binary case

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Given two pointed structures  $(\mathfrak{A}, a)$  and  $(\mathfrak{B}, b)$   
over a vocabulary which contains relation symbols of arity at most 2,

The pointed structure  $(\mathfrak{C}, c) = \text{Subst}((\mathfrak{A}, a), (\mathfrak{B}, b))$  is defined by:

- The universe of  $\mathfrak{C}$  is  $A \cup B - \{a\}$ .
- The context  $c$  is given by  $b$ , i.e.,  $c = b$ .
- For  $R \in \bar{R}$  of arity 1,  $R^{\mathfrak{C}}$  is defined by  $R^{\mathfrak{C}} = (R^{\mathfrak{A}} \cap (A - \{a\})) \cup R^{\mathfrak{B}}$ .
- For  $R \in \bar{R}$  of arity 2,  $R^{\mathfrak{C}}$  is defined by  $R^{\mathfrak{C}} = (R^{\mathfrak{A}} \cap (A - \{a\})^2) \cup R^{\mathfrak{B}} \cup I$   
where

$$\begin{aligned}
 I = & \{(a, x) \in A \times B : (a, b) \in R^{\mathfrak{A}}\} \\
 & \cup \{(x, a) \in B \times A : (b, a) \in R^{\mathfrak{A}}\} \\
 & \cup \{(x, x) \in B \times B : (a, a) \in R^{\mathfrak{A}}\}
 \end{aligned}$$

## Specker index, definition

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Let  $\mathcal{C}$  be a class of  $\tau = \bar{R}$ -structures, where  $\tau$  contains only relation symbols of arity at most 2.

We define an equivalence relation between  $\bar{R}$ -structures depending on  $\mathcal{C}$ .

(i)  $\mathfrak{A}_1$  and  $\mathfrak{A}_2$  are equivalent with respect to  $\mathcal{C}$ ,

$$\mathfrak{A}_1 \sim_{Su(\mathcal{C})} \mathfrak{A}_2$$

if for every pointed structure  $(\mathfrak{G}, s)$  we have that

$$Subst((\mathfrak{G}, s), \mathfrak{A}_1) \in \mathcal{C} \text{ iff } Subst((\mathfrak{G}, s), \mathfrak{A}_2) \in \mathcal{C}$$

(ii) The *Specker index* of  $\mathcal{C}$  is the number of equivalence classes of  $\sim_{Su(\mathcal{C})}$ .



## Specker index, properties

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### Proposition 10 (Specker)

Let  $\mathcal{C}$  be a class of  $\tau = \bar{R}$ -structures, where  $\tau$  contains only relation symbols of arity at most 2.

- (i) If  $\mathcal{C}$  is *MSOL*-definable then  $\mathcal{C}$  has finite Specker index.
- (ii) There are continuum many classes of finite Specker index.  
In fact,  $\text{Cycle}(A)$  has Specker index at most 5.

If we allow additionally modular counting quantifiers (i.e use the logic *CMSOL*), definability still implies finite Specker index.

The class of Eulerian graphs is *CMSOL*-definable but not *MSOL*-definable.

## Substitution of structures, general case

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We can analyse substitution in a more general context:

Given two pointed relational  $\tau$ -structures  $(\mathfrak{A}, a)$  and  $(\mathfrak{B}, b)$  without restrictions on the arity of the relation symbols in  $\tau$ .

We form a new pointed structure

$$(\mathfrak{C}, c) = \text{Subst}((\mathfrak{A}, a), (\mathfrak{B}, b))$$

defined as follows:

- The universe of  $\mathfrak{C}$  is  $A \cup B - \{a\}$ .
- The context  $c$  is given by  $b$ , i.e.,  $c = b$ .
- For  $R \in \bar{R}$  of arity  $r$ ,  $R^{\mathfrak{C}}$  is defined by

$$R^{\mathfrak{C}} = (R^{\mathfrak{A}} \cap (A - \{a\})^r) \cup R^{\mathfrak{B}} \cup I$$

where for every tuple in  $R^{\mathfrak{A}}$  which contains  $a$ ,  $I$  contains all possibilities of tuples obtained by replacing these occurrences of  $a$  with (identical or differing) members of  $B$ .

Analyze the properties of this new substitution and its corresponding Specker index.

You can try variations on the definition of  $I$ .

## Graphs of unbounded degree, II

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### Theorem 11 (Specker and Blatter, 1981)

Let  $\mathcal{C}$  be a class of  $\tau = \bar{R}$ -structures, such that

- (i)  $\tau$  contains only relation symbols of arity at most 2.
- (ii)  $\mathcal{C}$  has finite Specker index.

Then  $d_{\mathcal{C}}$  satisfies a linear modular recurrence relation.

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The full proof of the Specker-Blatter Theorem is left as project.

So is Fischer's counterexample of arity 4.

## The Specker-Blatter Theorem and Hamiltonian graphs

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Let  $HAM$  be the class of graphs with universe  $[n]$  which do have a Hamiltonian cycle, and let  $HAM(d)$  be those graphs in  $HAM$  which have degree at most  $d$ .

- $Ham$  has infinite Specker index, hence is not  $MSOL$ -definable.

Use  $Star_n$  with center  $a$  and note that  $Subst(Star_m, a, E_n) = K_{m,n}$ , which is hamiltonian iff  $m = n$ .

Note for non-definability we used a transduction mapping  $a^n b^m$  on words to  $K_{m,n}$ , too.

- Is  $HAM(d)$   $MSOL$ -definable?
- As  $HAM(d)$  is connected its  $DU$ -index is 2. What is its Specker index?
- Do  $d_{HAM}$  and  $d_{HAM(d)}$  satisfy linear modular recurrences?
- Try to compute  $d_{HAM}$  and  $d_{HAM(d)}$ .

## Research problems stemming from the Specker-Blatter Theorem

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There are many unresolved questions here, suitable for research.

- Is the Specker-Blatter Theorem true for arity 3?
- Can we strengthen the assumption on finite Specker index such that the Specker-Blatter Theorem would work for arbitrary finite relational vocabularies?
- *subst* is an *MSOL*-smooth operation on vocabularies with relation symbols of arity at most 2. The same holds for *fuse*.  
Explore the relationship between the operation *subst* and *fuse* and the general notion of *patch*-width.
- Analyze the relationship between *recognizable* classes of structures in the sense of Courcelle, and finite Specker index (and related notions).