

Logical Methods in Combinatorics 236605 (2009/10)

Homework to Lecture 1 and 2

- (i) Topologies and preorders.
- (ii) Counting 1-regular graphs (= perfect matchings).
- (iii) Check monotonicity, hereditariness and definability of selected graph properties.
- (iv) Forbidden (induced) subgraphs
- (v) **MSOL**-definability of regular languages
- (vi) 3-colorability

Topologies and preorders, I

A *preorder* R is a transitive and reflexive binary relation.

If additionally $R(a, b) \wedge R(b, a) \Rightarrow a = b$, R is a *partial order*.

$PreO_n$ is the number of preorders on the set $\{1, \dots, n\}$.

Theorem:

The number of preorder on a finite set equals the number of topologies, i.e., $PreO_n = T_n$.

Complete and verify the following proof sketch.

Topologies and preorders, II

Proof: (From topology to preorder)

Let $S = \{1, \dots, n\}$.

Let $Cl : P(S) \rightarrow P(S)$ be any function. closure operation coming from some topology.

Define $cl : S \rightarrow P(S)$ by $cl(a) = Cl(\{a\})$.

Define $R_{Cl}(a, b)$ by $a \in cl(b)$.

Claim: R_{Cl} is a preorder iff Cl is the closure operation from some topology.

This uses that S is finite.

Topologies and preorders, III

Proof: (From preorder to topology)

Let R be a preorder on S .

We define a function $cl_R : S \rightarrow P(S)$ by

$$cl_R(b) = \{a \in S : R(a, b)\}.$$

As for finite sets, the topology is defined, once we know the closures of the singletons, we define for $X \subseteq S$

$$Cl(X) = \bigcup_{a \in X} cl(a)$$

To verify that this works we use the previous claim .

Counting perfect matchings

Let \mathcal{P} consist of graphs which are 1-regular
(disjoint union of non-connected edges, perfect matchings).

- (i) \mathcal{P} is not hereditary. **Verify!**
- (ii) \mathcal{P} is *FOL*-definable. **Write the formula explicitly!**
- (iii) **Compute and verify the density** $f_{\mathcal{P}}(m)$. and $sp(\mathcal{P})$.
- (iv) **Compute and verify the fractions and probabilities** $\mu_m(\mathcal{P})$ and $\mu(\mathcal{P})$.

Monotonicity and Definability

Check for monotonicity, hereditariness and definability in **FOL**, **MSOL** and **SOL**:

- (i) Connected graphs, graphs of diameter d ;
- (ii) regular graphs, r -regular graphs;
- (iii) Graphs with an Eulerian or Hamiltonian path (circuit);
- (iv) Planar graphs,
- (v) bipartite graphs, r -colorable graphs.

Forbidden subgraphs

Choose your favorite graph H .

Consider the classes of graphs $Forb(H)$ and $iForb(H)$ of graphs with no subgraph, resp. induced subgraph isomorphic to H .

- For which H are these classes monotone, hereditary?
Try to formulate a general answer based on properties of H .
- Show that for each H both are FOL -definable.
- Try to compute the density function and asymptotic probabilities for various choices of H .
- Find the optimal results in the literature.

MSOL on words (Homework)

Theorem:[Büchi, Elgot, Trakhtenbrot, 1961]

A class of binary words is:
recognizable by a finite
(non-deterministic) automaton
iff it is **MSOL**-definable
(iff it is regular).

Example: $(101 \vee 1001)^*$
101 1001 101 101 1001 1001 101.....

Exercise: Find the **MSOL**-formula.

k -colorability

Homework:

- (i) Compute the speed of $3COL$.
- (ii) Is $3COL$ of the form $Forb_{ind}(H)$?
- (iii) Is $3COL$ **FOL**-definable?
- (iv) Define k -colorability and show, that for fixed k , it is **MSOL**-definable.