

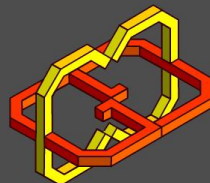


# Geometric Covering

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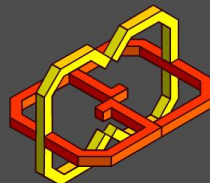
MSc Thesis



# Introduction

Geometric Covering (GC) queries appear in numerous applications:

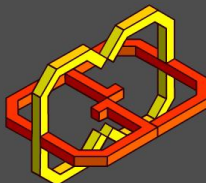
- Mold design in manufacturing
- Inspection
- Security and surveillance
- Placements of cellular antennas
- Illumination design
- Spraying of paint



# Layout of the Rest of the Talk

We are focusing on mold-design and security.

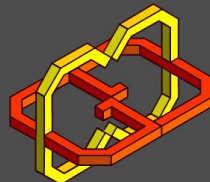
- ❑ Related work in mold-design and security.
- ❑ A generic unified framework for answering geometric covering.
- ❑ Geometric Covering is an *NP*-hard problem.
- ❑ Examples of the generic framework as implemented in a 3D mold-design and security.
- ❑ Conclusions and future work.



# Related Work I

## Mold design

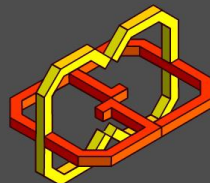
- ❑ 2-pieces-mold polygonal decomposition in  $R^3$   
[Ahn02, Khardekar06, Chen06]
- ❑  $n$ -pieces-mold polygonal decomposition in  $R^3$  [Liu09, Priyadarshi04, Stoyan10]
- ❑ 2-pieces-mold freeform surface decomposition in  $R^3$   
[Elber04]
- ❑ Algebraic analysis of visibility of freeforms in  $R^3$   
[Seong06]
- ❑ **Nothing** so far on automatic  $n$ -pieces-mold freeform decomposition in  $R^3$



# Related Work II

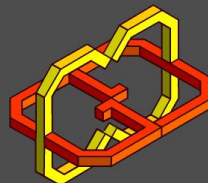
## Security

- ❑ Polygonal 2.5D terrain where  $z = f(x, y)$ .
- ❑ Guards on the vertices or above them [Lee91, Goodchild89]
- ❑ Edge guards [Bose96, Bose97]
- ❑ Different greedy solutions [Goodchild89, Kaucic04]
- ❑ Guards limited to strategic locations [Kim04]
- ❑ Calculating partial visibility [Franklin94, Rana03]



# Set-Cover I

- ❑ Set-cover (SC) is a classic computer science query.
- ❑ SC is considered a very hard problem to solve (*NP* hard).
- ❑ Given some universe  $U$  and a family  $F$  of subsets of  $U$  which their union equals  $U$ , a cover of  $U$  is a subfamily of  $F$  whose union still equals  $U$ .
- ❑ In SC we are seeking a cover with minimal number of subsets.

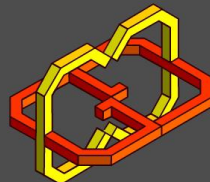
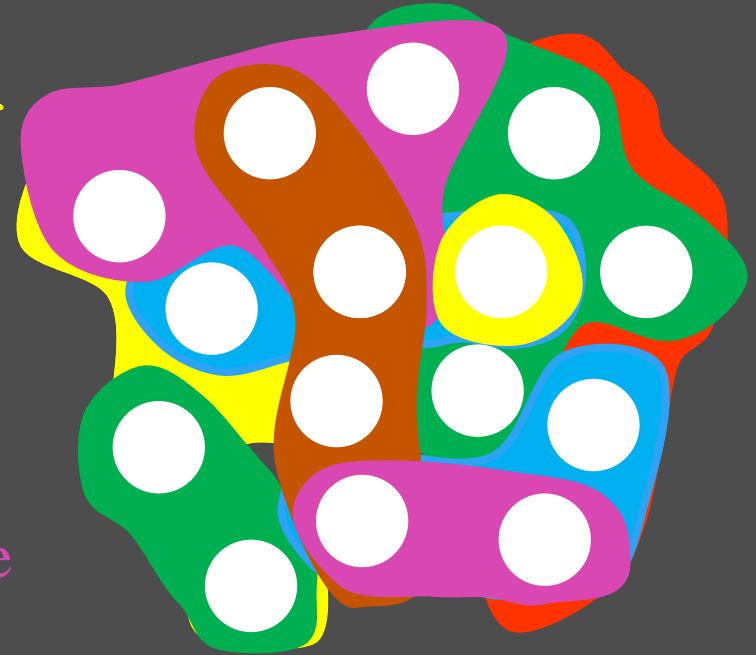


# Set-Cover II

- ❑ The universe  $U$  is a set of circles.
- ❑ A subset of  $U$  is a group of circles.
- ❑ The family  $F$  is all these groups of circles.
- ❑ The subfamily  $F_1$  is the brown, yellow, blue and green groups.  
 $F_1$  is a cover of  $U$ .
- ❑ The subfamily  $F_2$  is the red, purple and yellow group.

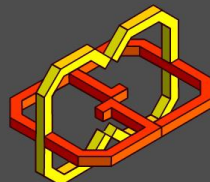
$F_1$  is a minimal cover of  $U$ .

We will now show a reduction from  
GC problems to SC problems.



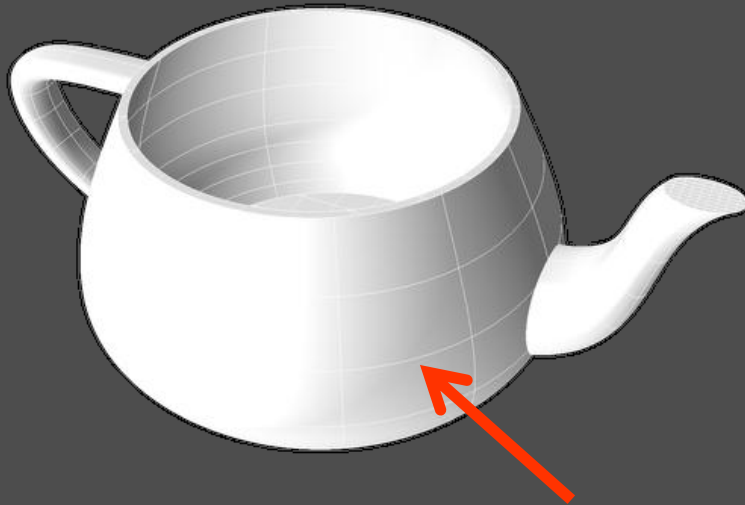
# Visibility Map I

- ❑ We receive a 2 manifold geometry in  $R^3$ ,  $C$ , which has a parameterization  $x_{uv}$ ,  $y_{uv}$ ,  $z_{uv}$ .
- ❑ The domain  $D_C$  of  $C$  is a 2-dimensional box, a rectangle, possibly trimmed.
- ❑ We are creating a discrete representation of  $D_C$  as an image, as a visibility map.
- ❑ The visibility map can serve as a controlled approximation for the coverage of  $C$ .





# Visibility Map II

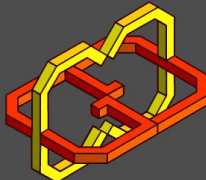


The Utah Teapot with its interior curved in.

The visibility map of the **outer** body of the Teapot



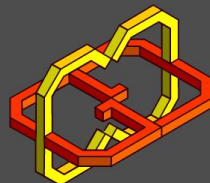
- Visible locations are set to **white**.
- Hidden locations are set to **black**.
- Trimmed away bits are set to **green** - don't care.



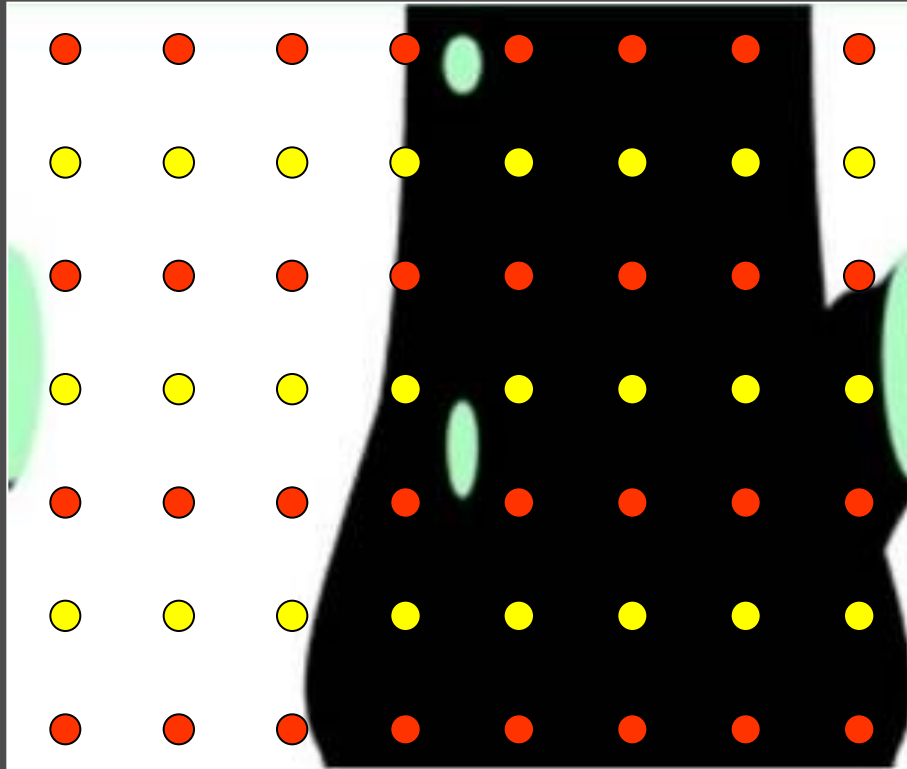
# Visibility Map III

Linearize the visibility map, as a vector of bits as follow:

- ❑ Don't care locations are simply skipped.
- ❑ Each bit is either 1 (visible pixel) or 0 (hidden pixel).
- ❑ Sequence the 1/0 bits in some order over the visibility map (for example: left to right, top to bottom).



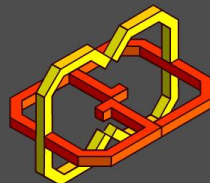
# Visibility Map IV



Visibility map of  
 $8 \times 7$

1111000111100001111000011110000011100000111000001110000011100000

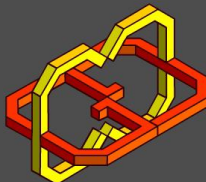
Vector of 56 bits



# Set-Cover II

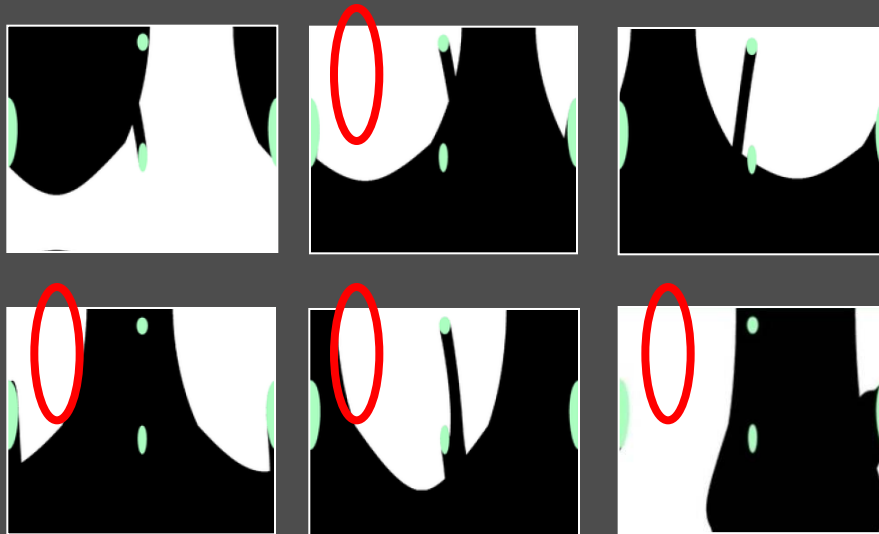
Set-cover can be clearly applied to vectors of bits:

- ❑ The universe  $U$  is the domain  $D_C$ .
- ❑ A subset of  $U$  is a vector of bits.
- ❑ A family  $F$  of subsets of  $U$  is a set of vectors of bits from different views around the geometry  $C$ .
- ❑ A cover of  $U$  is a subfamily of  $F$ , a set of vectors of bits which their union equals  $D_C$ .

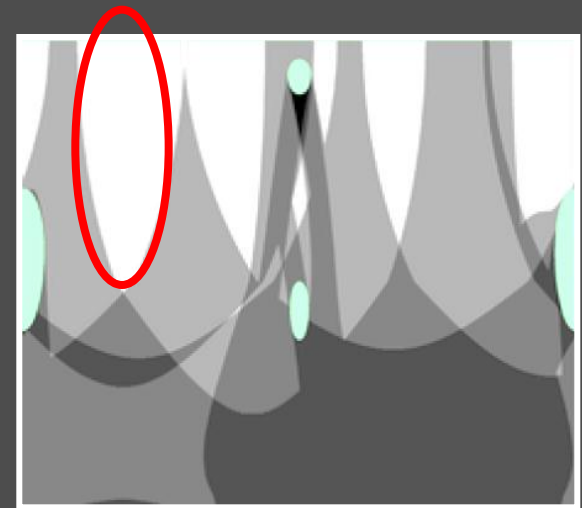


# Set-Cover III

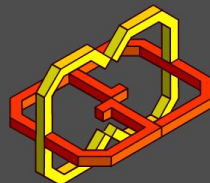
Subfamily of the set of  
visibility maps



The union of the  
visibility maps

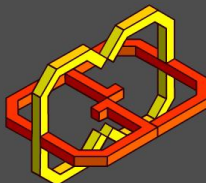


- The set-cover is done in the parametric domain.

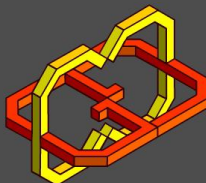
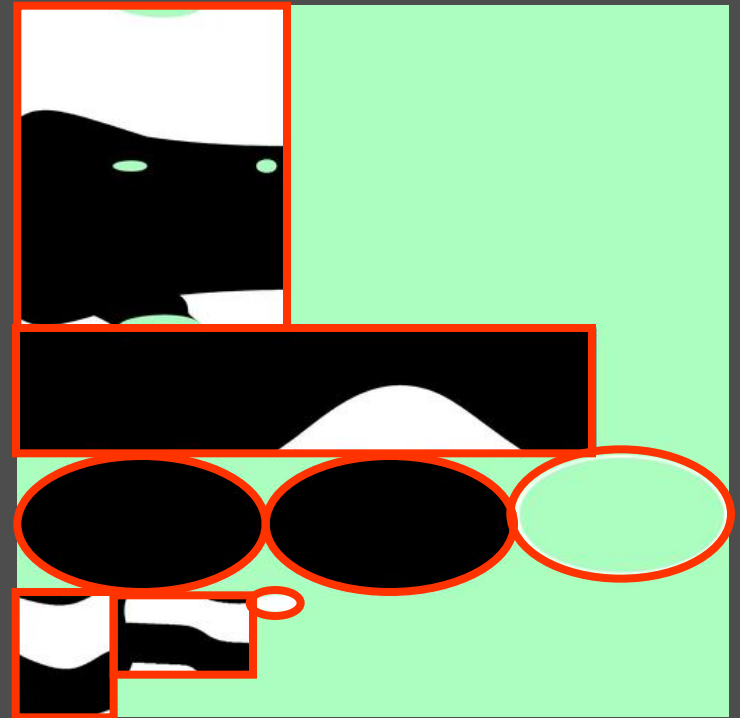
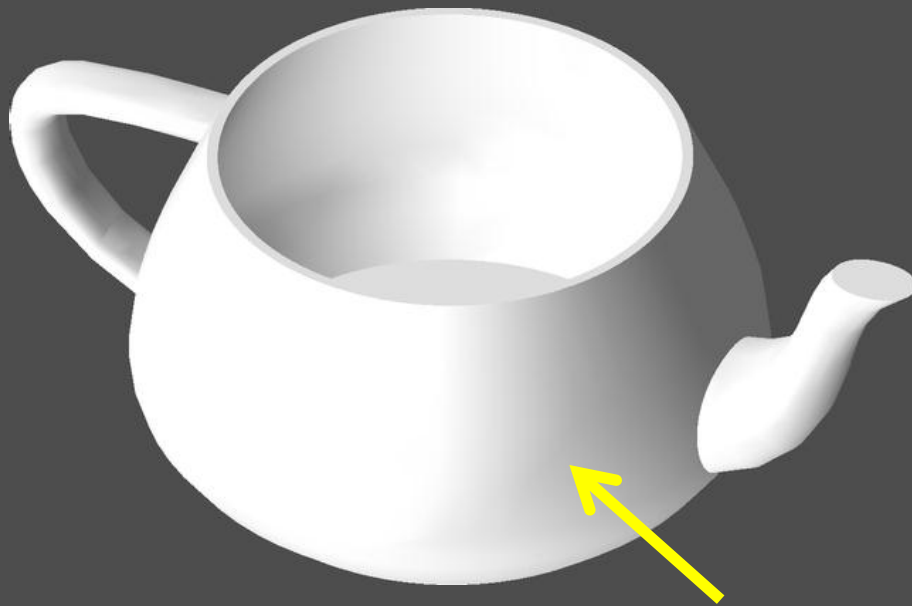


# Creating Visibility Maps I

- ❑ Input geometry  $C$  can be a surface or a set of surfaces, possibly trimmed.
- ❑ Each surface has its own rectangular domain, created independently of the other surfaces.
- ❑ We rearrange the domains of all the surfaces in one large image: The visibility map of  $C$ .

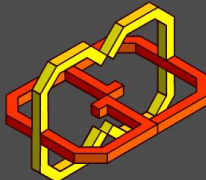


# Creating Visibility Maps II



# Creating Visibility Maps III

- Given  $C$  and  $D_C$ , the visibility map from direction  $V_i$  is computed as follow:
- The surface is tessellated into triangles.
- Two-rendering passes:
  - I. A regular (Z-buffer) rendering of  $C$  from  $V_i$  keeping only the Z-depth information, in  $ZBuffer(x, y)$ .
  - II. Scan conversion of  $C$  in the domain,  $D_C$ , and deciding visibility by comparing the Z-depths





# Creating Visibility Maps IV

## Pass II

A tessellation  $T = \{T_i\}$  of triangles with  $UV$  parametric coordinates is given.

**For** each triangle  $T_i$  in  $T$ , scan convert  $T_i$  by its  $UV$  coordinates.

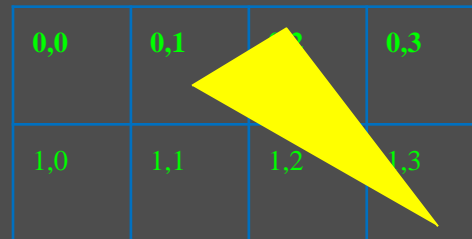
**For** each pixel  $p_{uv}$  in  $T_i$

$x_{uv}, y_{uv}, z_{uv} \leftarrow XYZ$  coordinates of  $p_{uv}$ ;

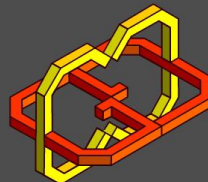
$VisMap(u, v) \leftarrow z_{uv} \approx ZBuffer(x_{uv}, y_{uv})$ ;

**EndFor**

**EndFor**



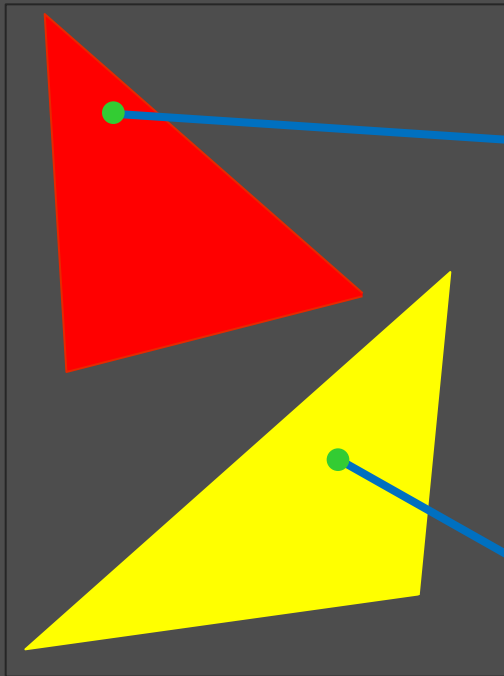
$UV$  Domain of  
 $4 \times 2$



# Creating Visibility Maps V

*UV* domain

pass II

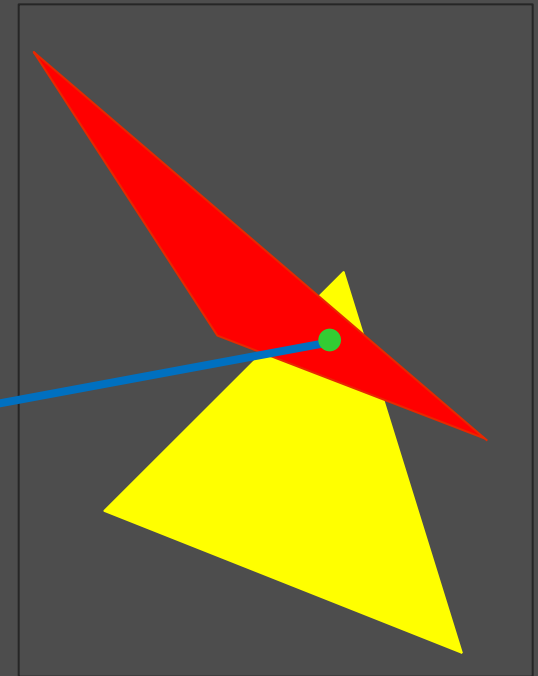


$(u_1, v_1)$   
 $(x, y, z_1)$

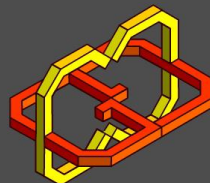
$(u_2, v_2)$   
 $(x, y, z_2)$

Euclidean space

pass I

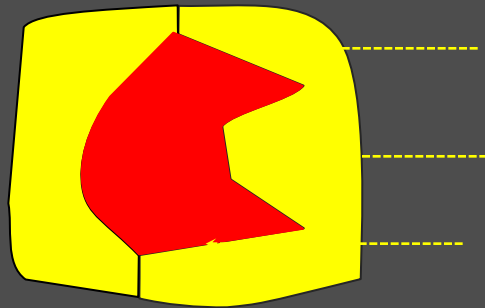


$$ZBuffer(x, y) \approx z_1$$



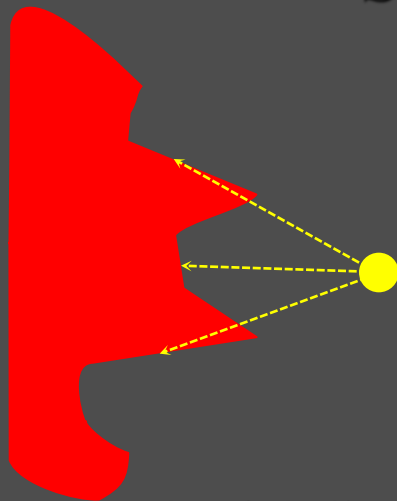
# Creating Visibility Maps VII

## Mold Design

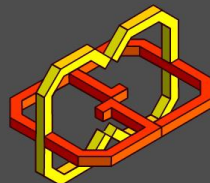


Orthographic  
projection

## Security

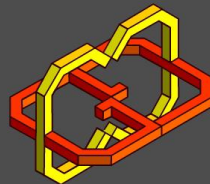
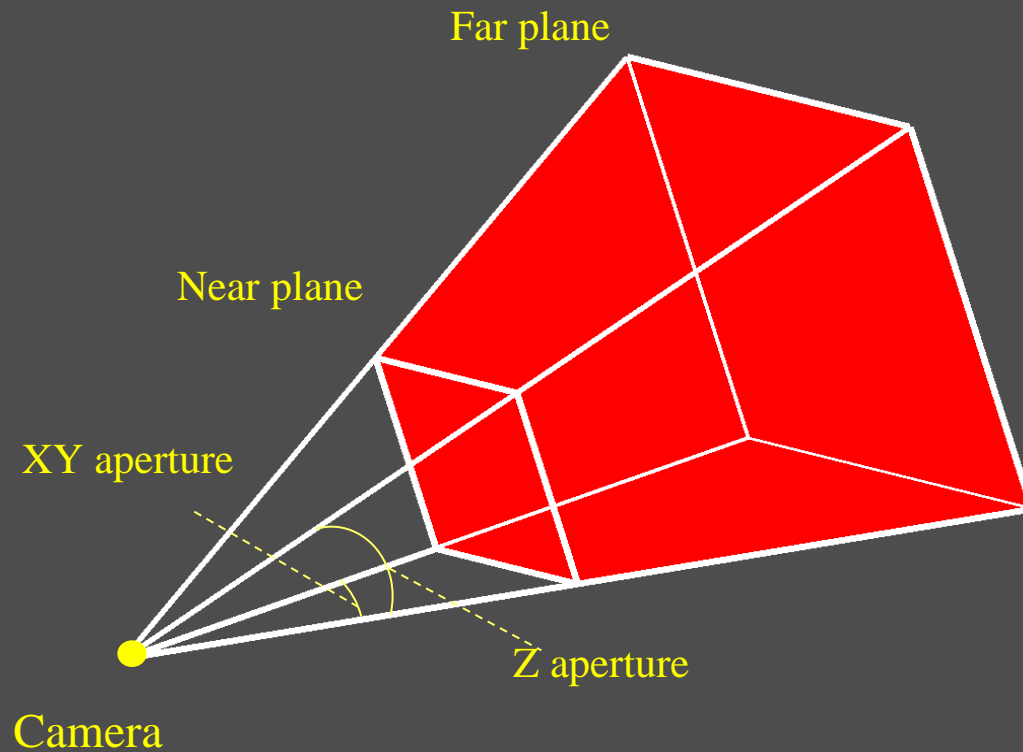


Perspective  
projection



# Creating Visibility Maps VIII

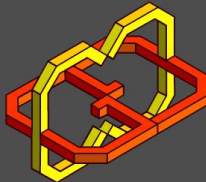
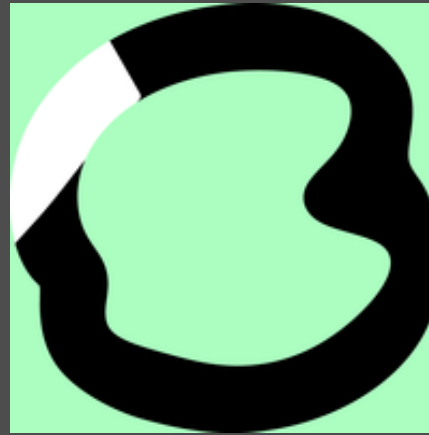
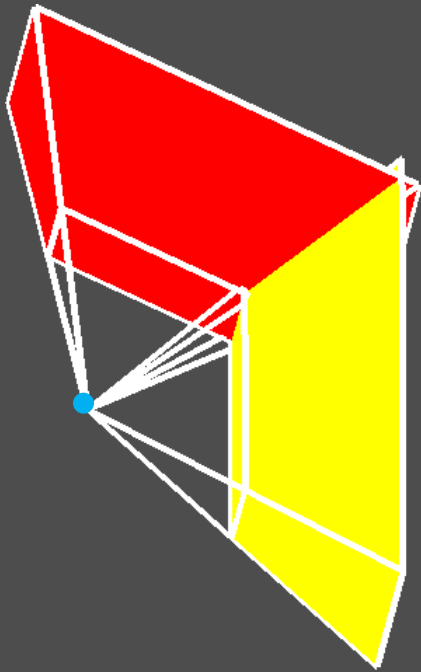
## Perspective projection I



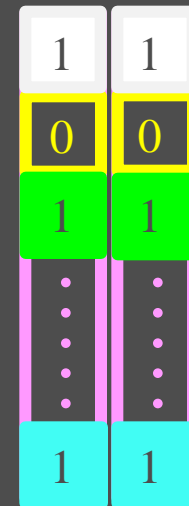
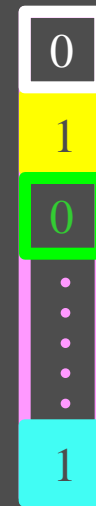
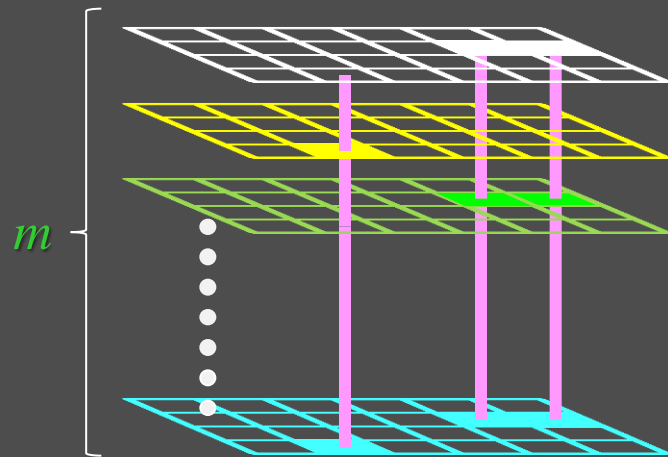
# Creating Visibility Maps IX

## Perspective projection II

Combining visibility maps

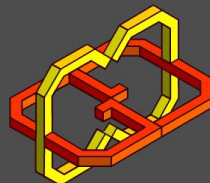


# Pixel Collapsing I



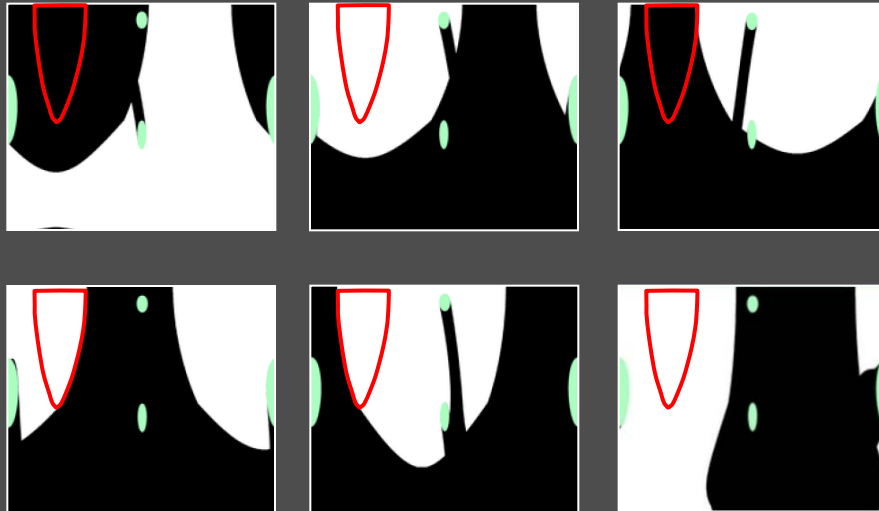
$n \times n \times m$

- ❑  $2^m$  possible pixels vector.
- ❑  $n^2$  different pixels vector at most.
- ❑ In practice, much less.

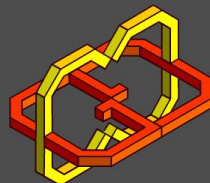
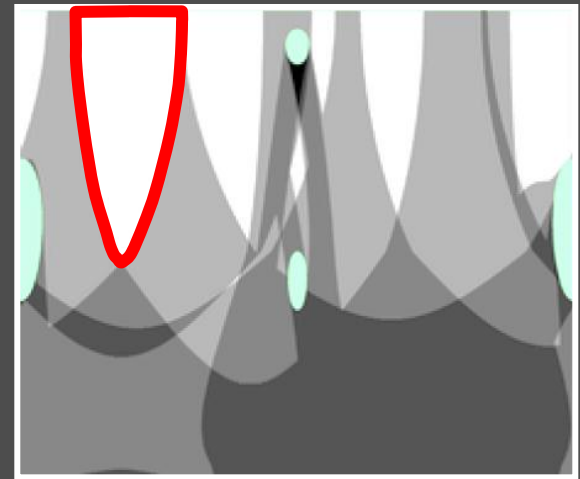


# Pixel Collapsing II

Subfamily of the set of  
visibility maps

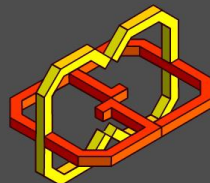


The union of the  
visibility maps



# Reduction from SC to GC I

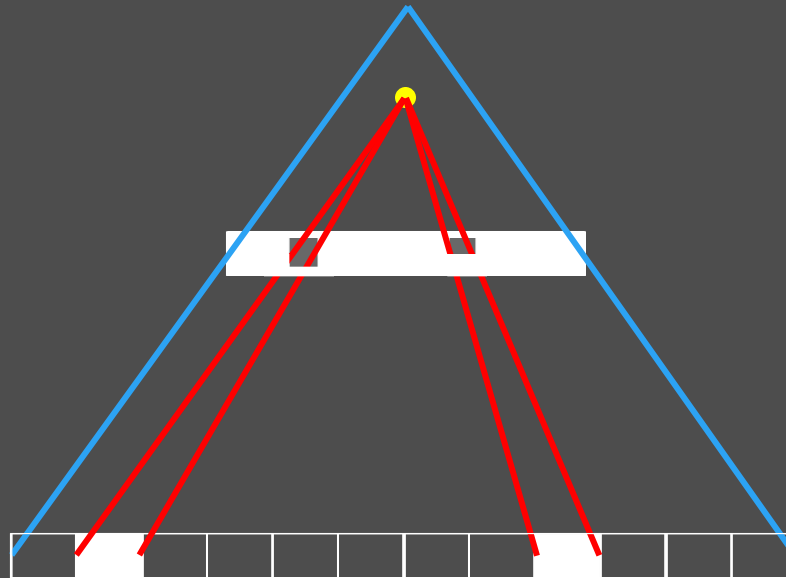
- ❑ We have shown a polynomial reduction from GC to SC. For completeness we will also show a polynomial reduction from SC to GC, proving that GC is *NP*-hard as SC is.
- ❑ We have a standard SC as described before.
- ❑ We will create a geometry corresponding to the universe  $U$ .
- ❑ We will create guards corresponding to the subsets of  $U$ .
- ❑ Solving the GC will solve the SC as well.





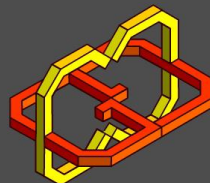
# Reduction from SC to GC II

Subset of  $U$  - a possible guard.



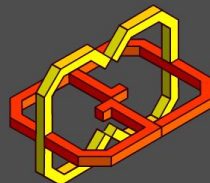
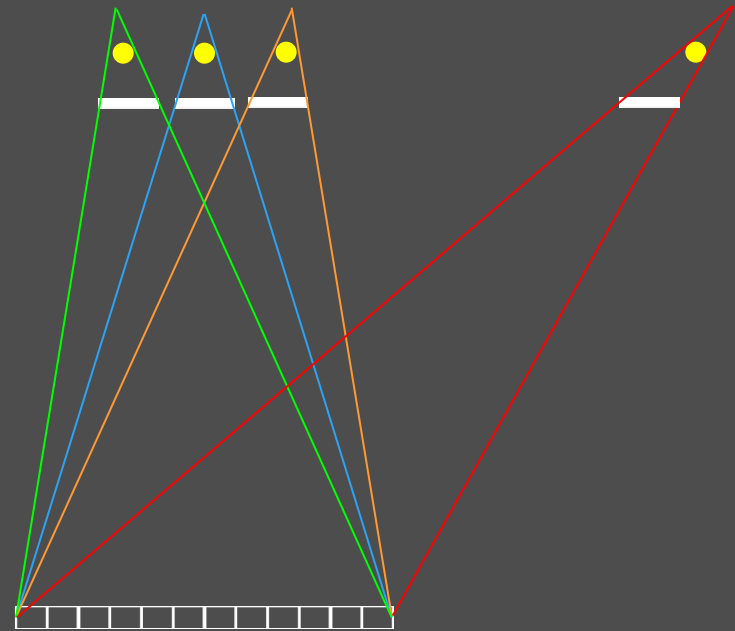
$U$  - a long strip.

Elements of  $U$  - regions on the strip.



# Reduction from SC to GC III

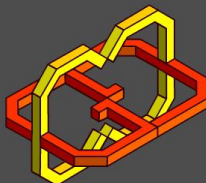
- $F$  - as many guards as are subsets in the problem, spread over the entire plane.
- All the upper strips are entirely covered by each of the guards.



# Examples

## General Notes

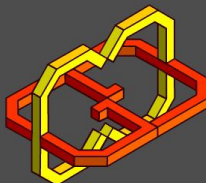
- ❑ The following examples were created using Visibility maps of size  $4096 \times 4096$ .
- ❑ Both exhaustive (exponential) set cover solution and greedy (non-optimal) solution were sought.
- ❑ All implementation is software based and with single thread.
- ❑ In the examples we seek high coverage percent rather than a complete coverage.



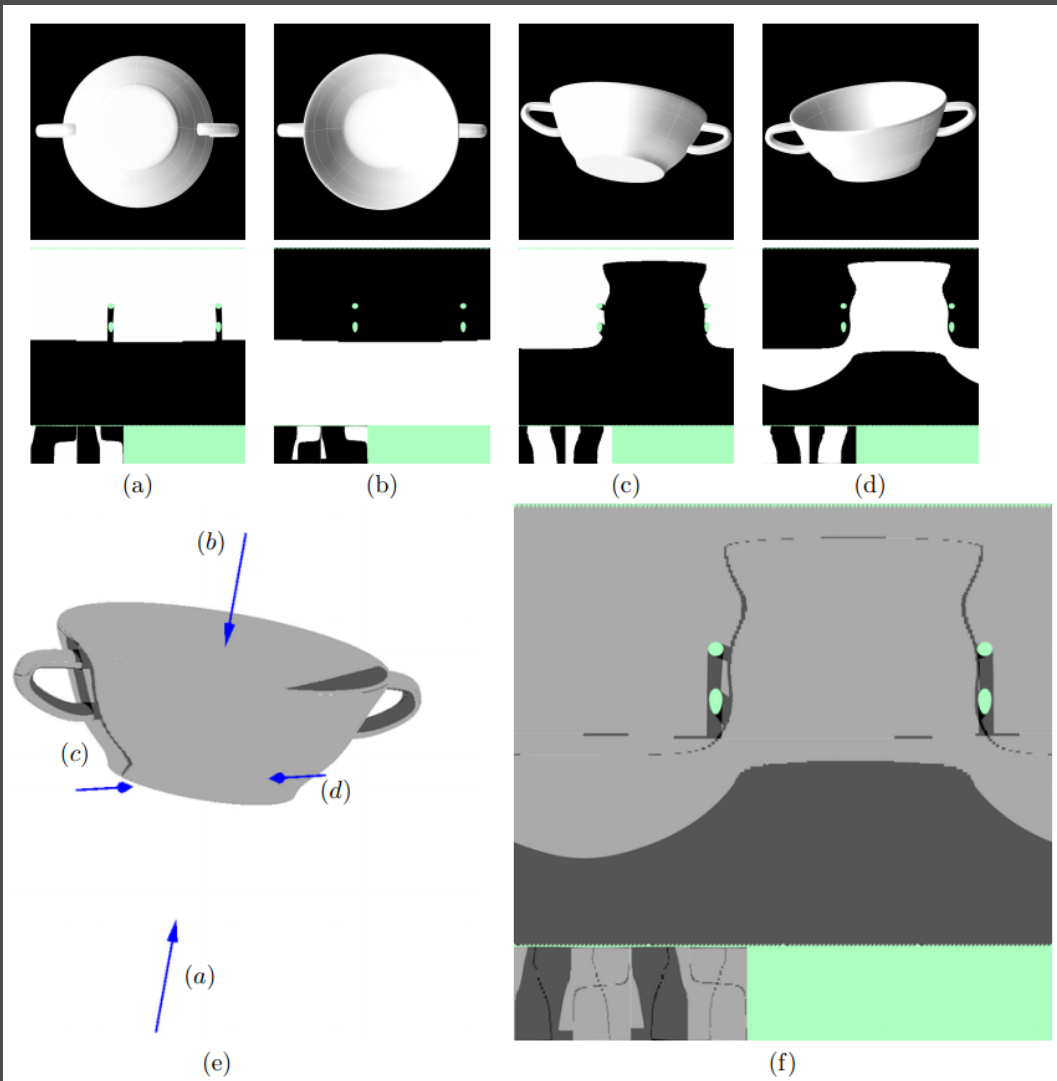
# Mold-Design Examples

## General Notes

- The following examples were created using 266 views:
  - 130 general views around  $S^2$ , duplicated as  $V$  and  $-V$ .
  - 6 views of  $\pm X$ ,  $\pm Y$ ,  $\pm Z$ .



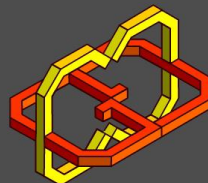
# Example – a Cup Model



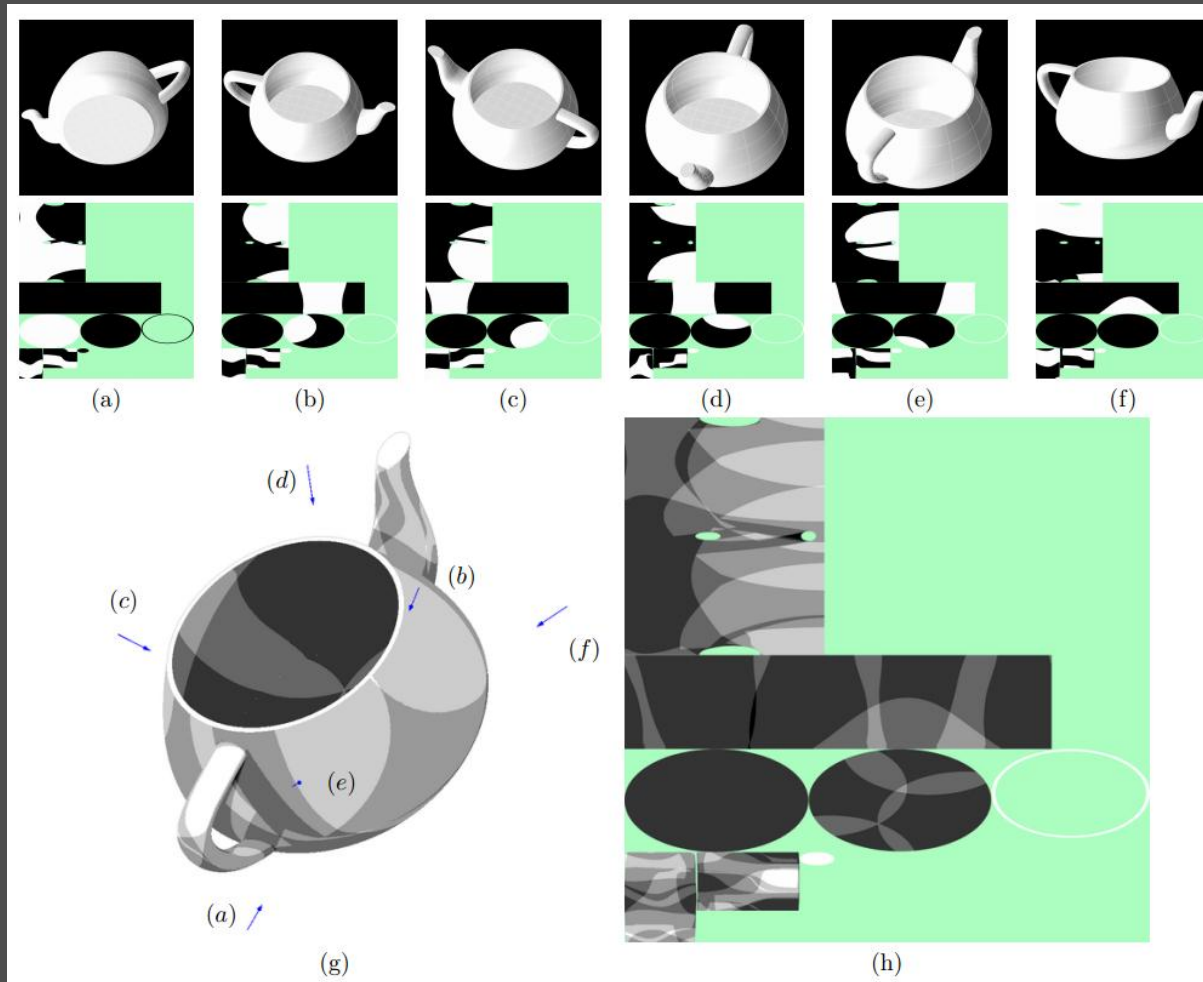
99.827% cover in greedy SC in ~4 seconds.

99.995% cover in exhaustive SC in ~10 hours.

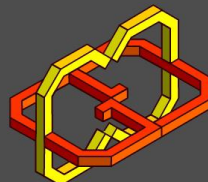
First two view directions 95% cover.



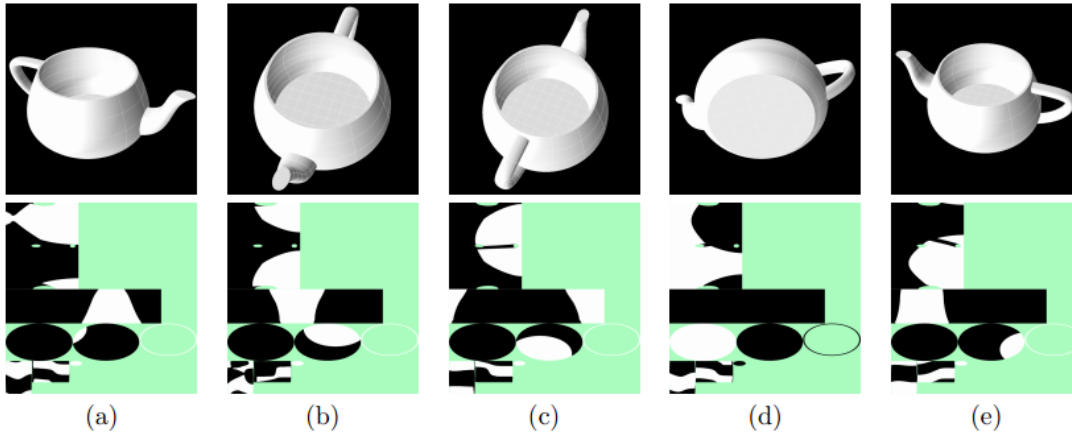
# Example – The Utah Teapot I



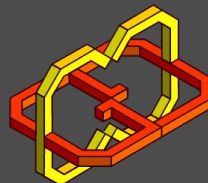
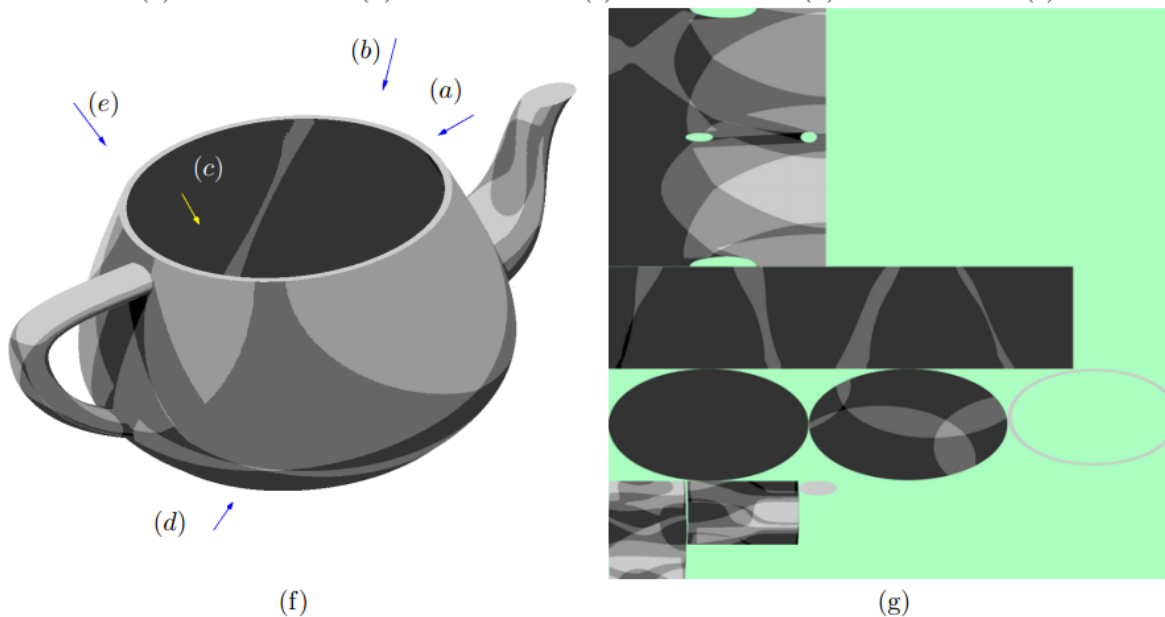
99.7% cover in greedy SC in ~6 seconds.



# Example – The Utah Teapot II



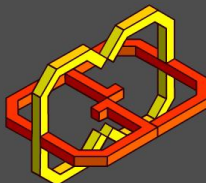
99.7% cover in  
exhaustive SC in ~433  
hours.



# Security Examples

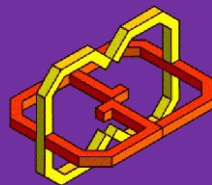
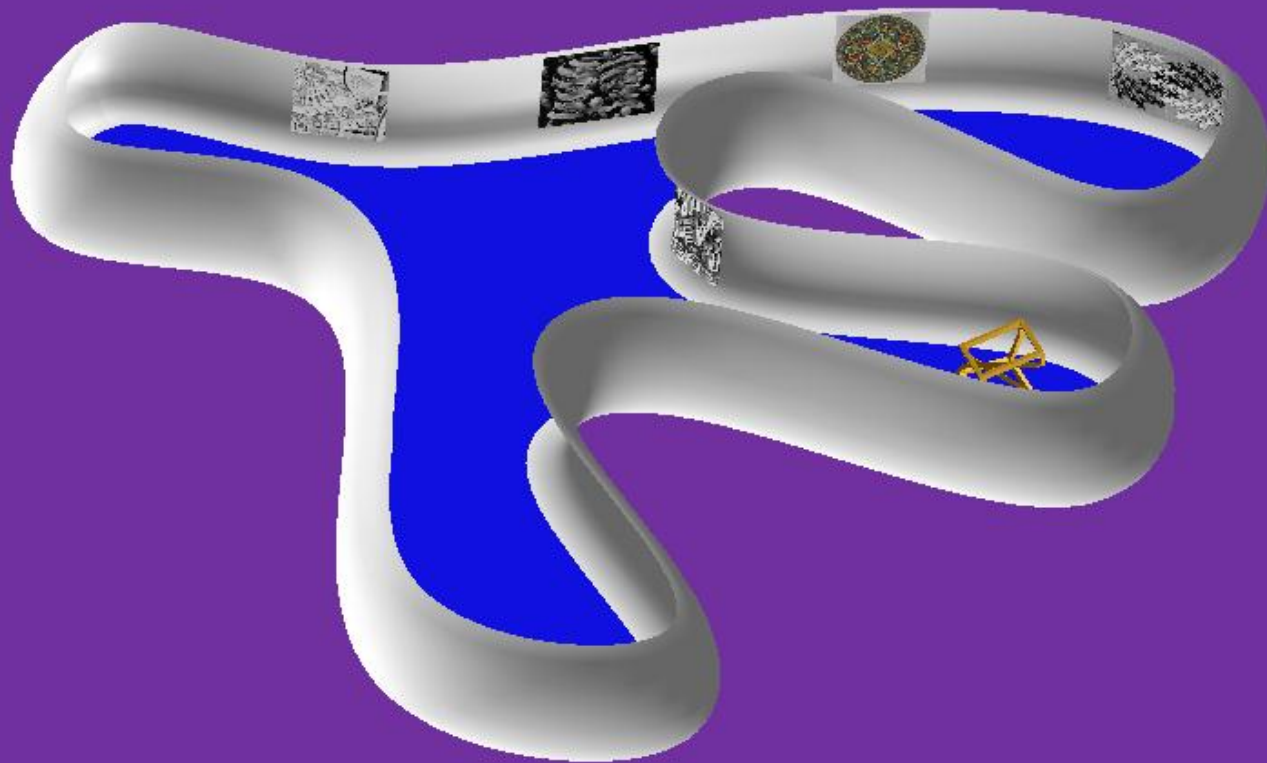
## General Notes

- ❑ The following examples were created using about 300 guards/cameras.
- ❑ The guards where evenly spread on a curve or a plane.

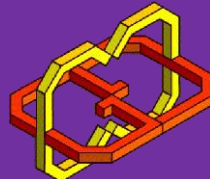
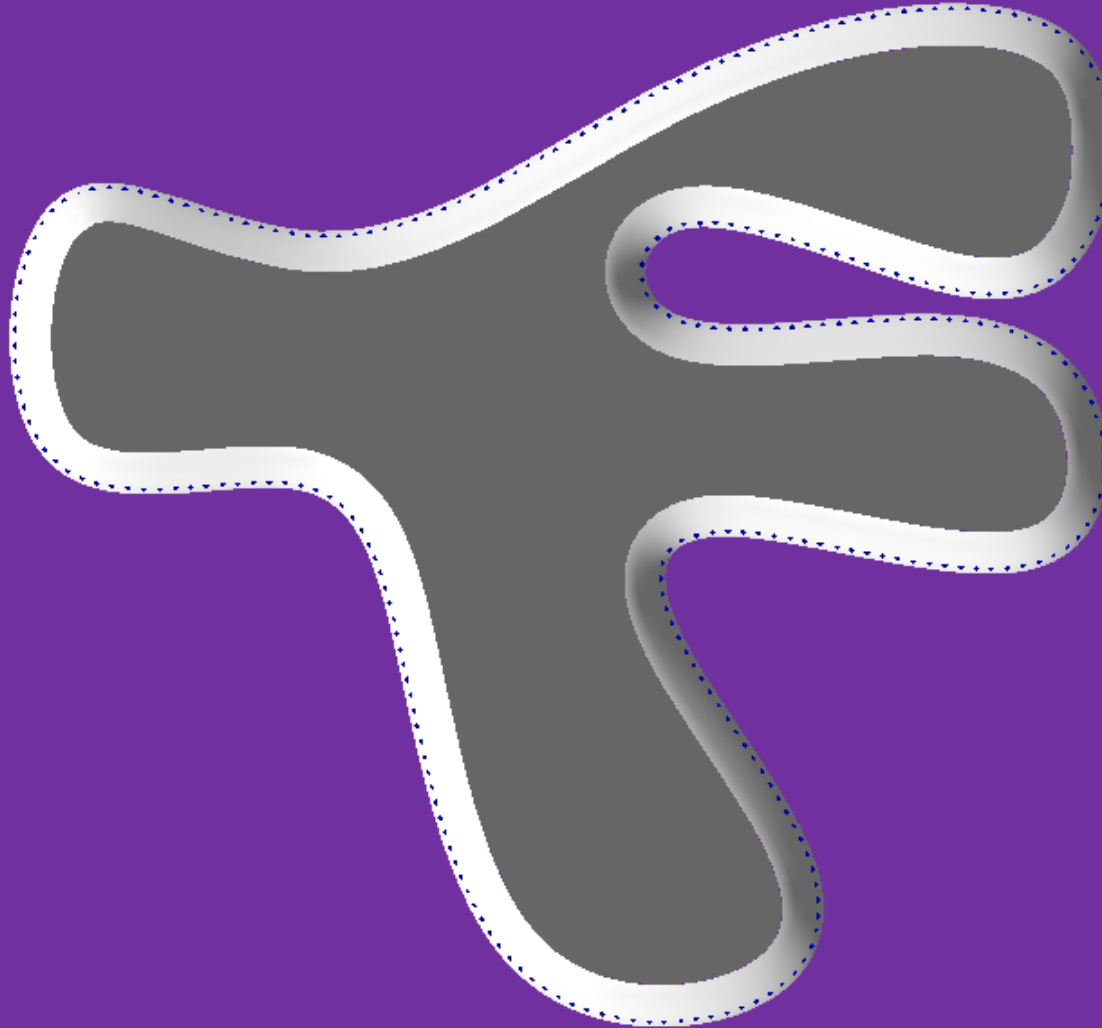




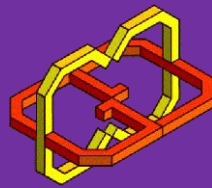
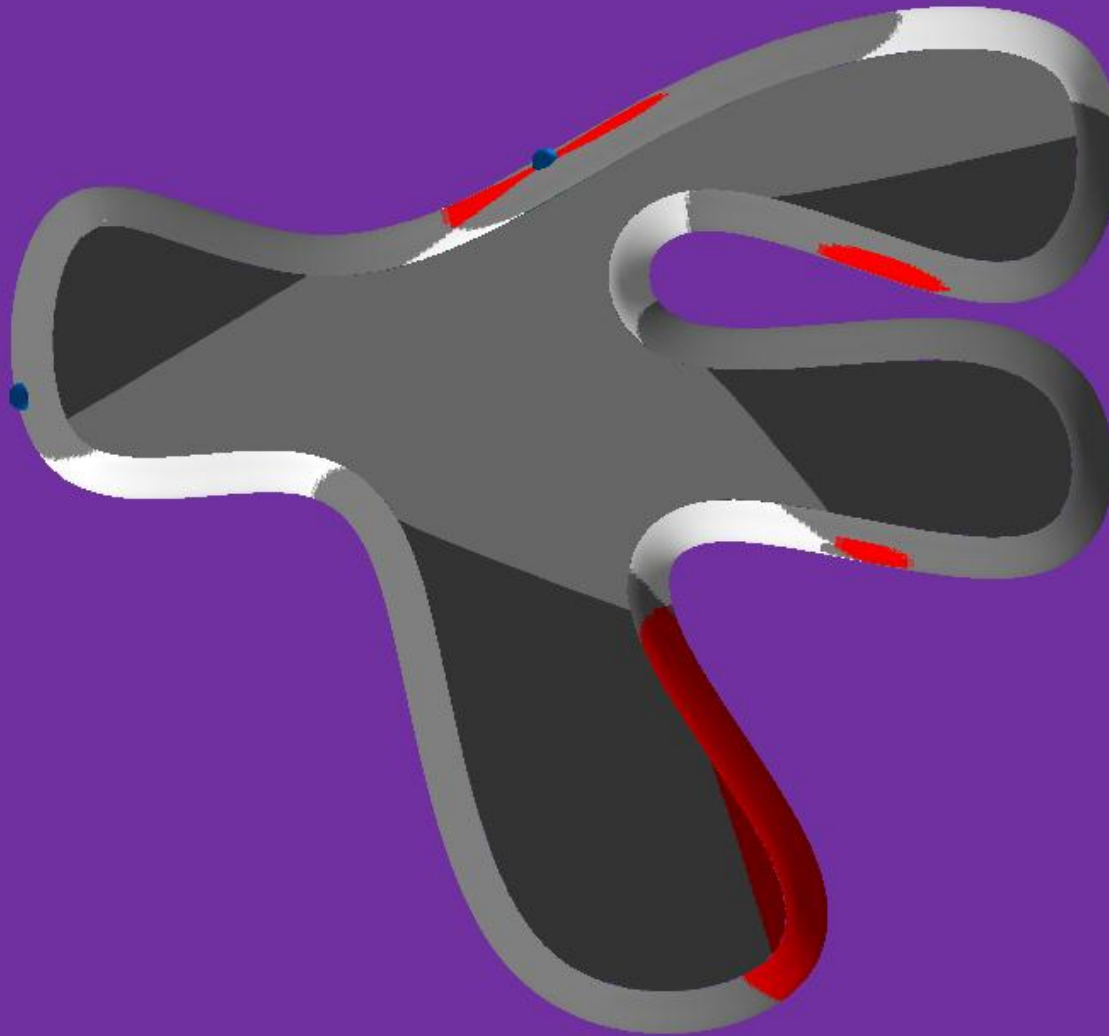
# A free form shape gallery



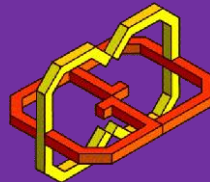
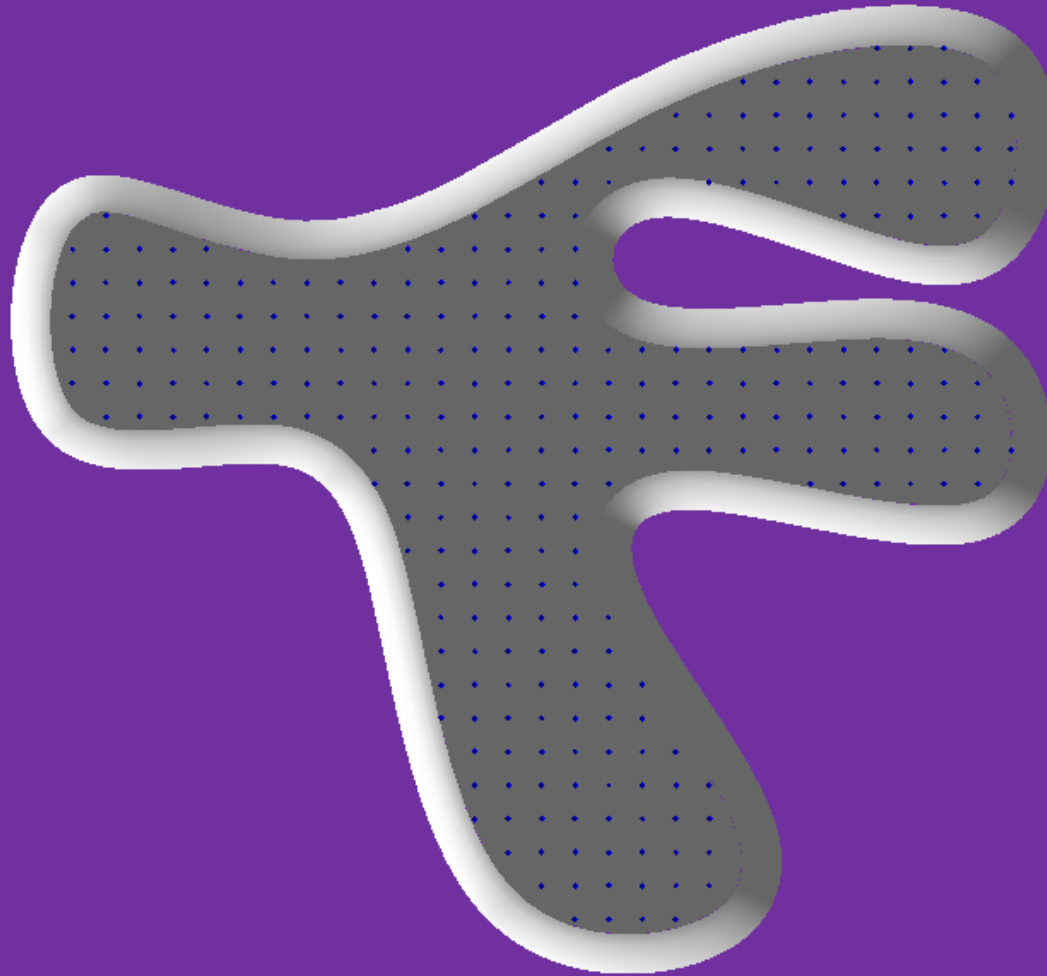
# Cameras on the walls



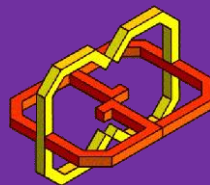
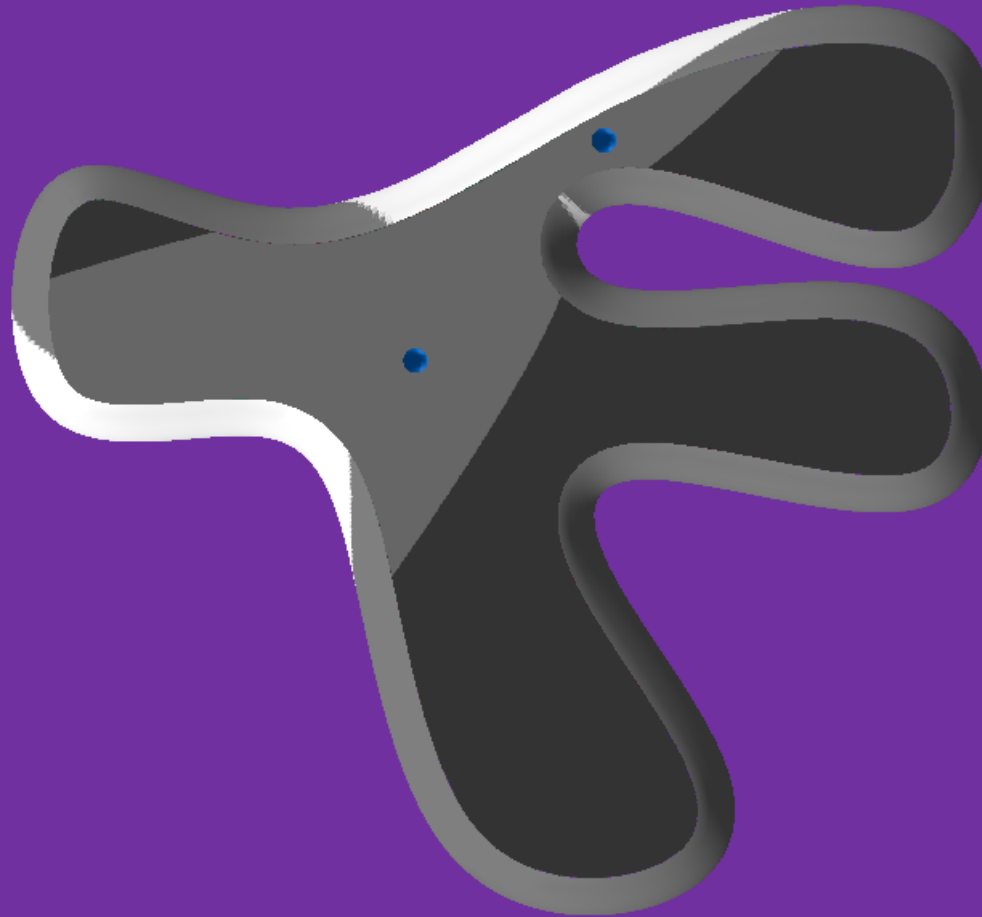
# Cameras on the wall - 2 cameras solution



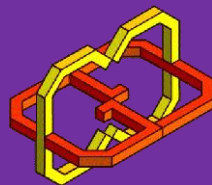
# Cameras on the ceiling



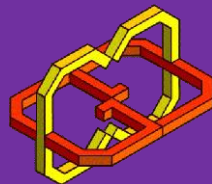
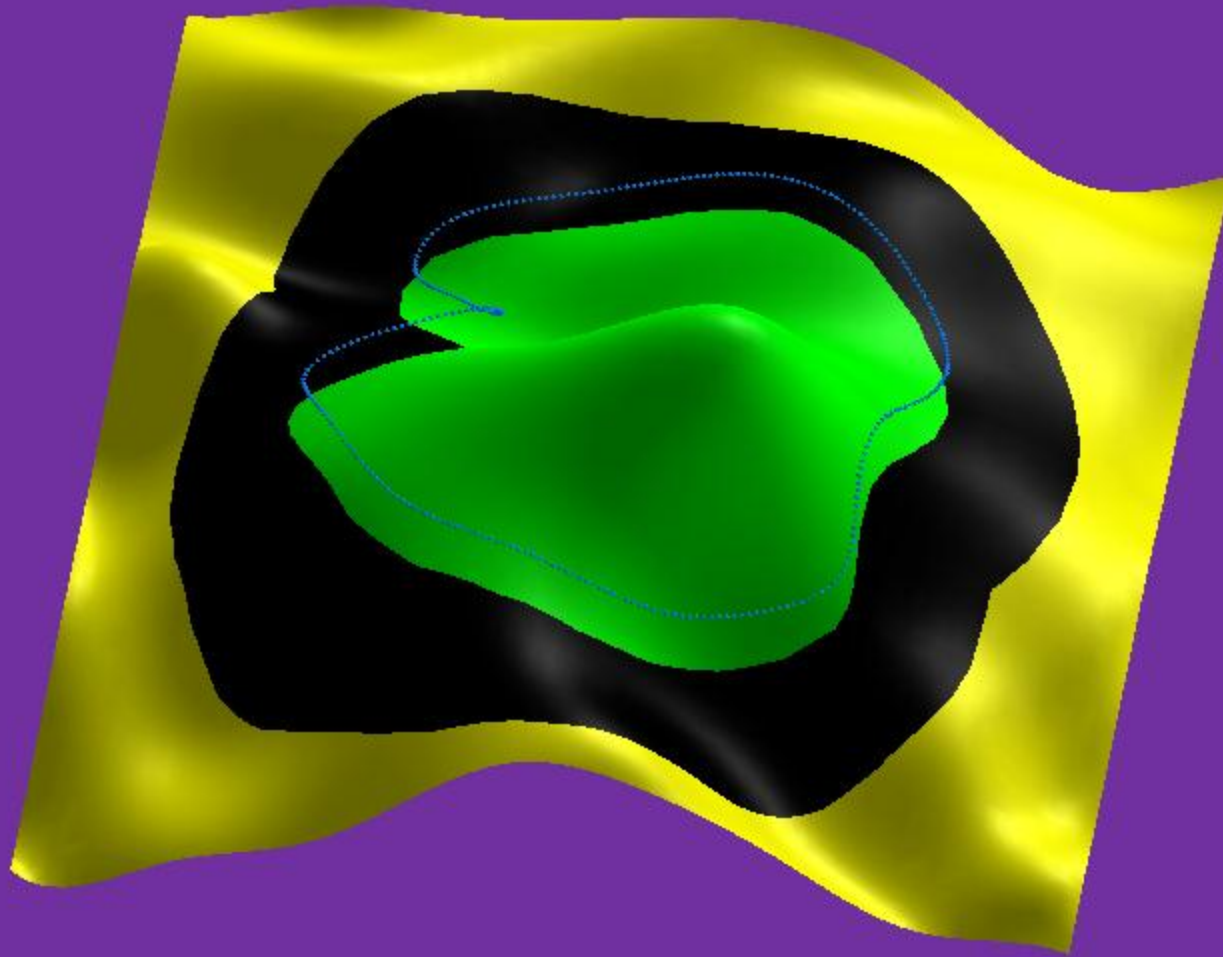
# Cameras on the ceiling - 2 cameras solution



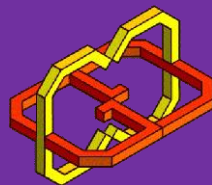
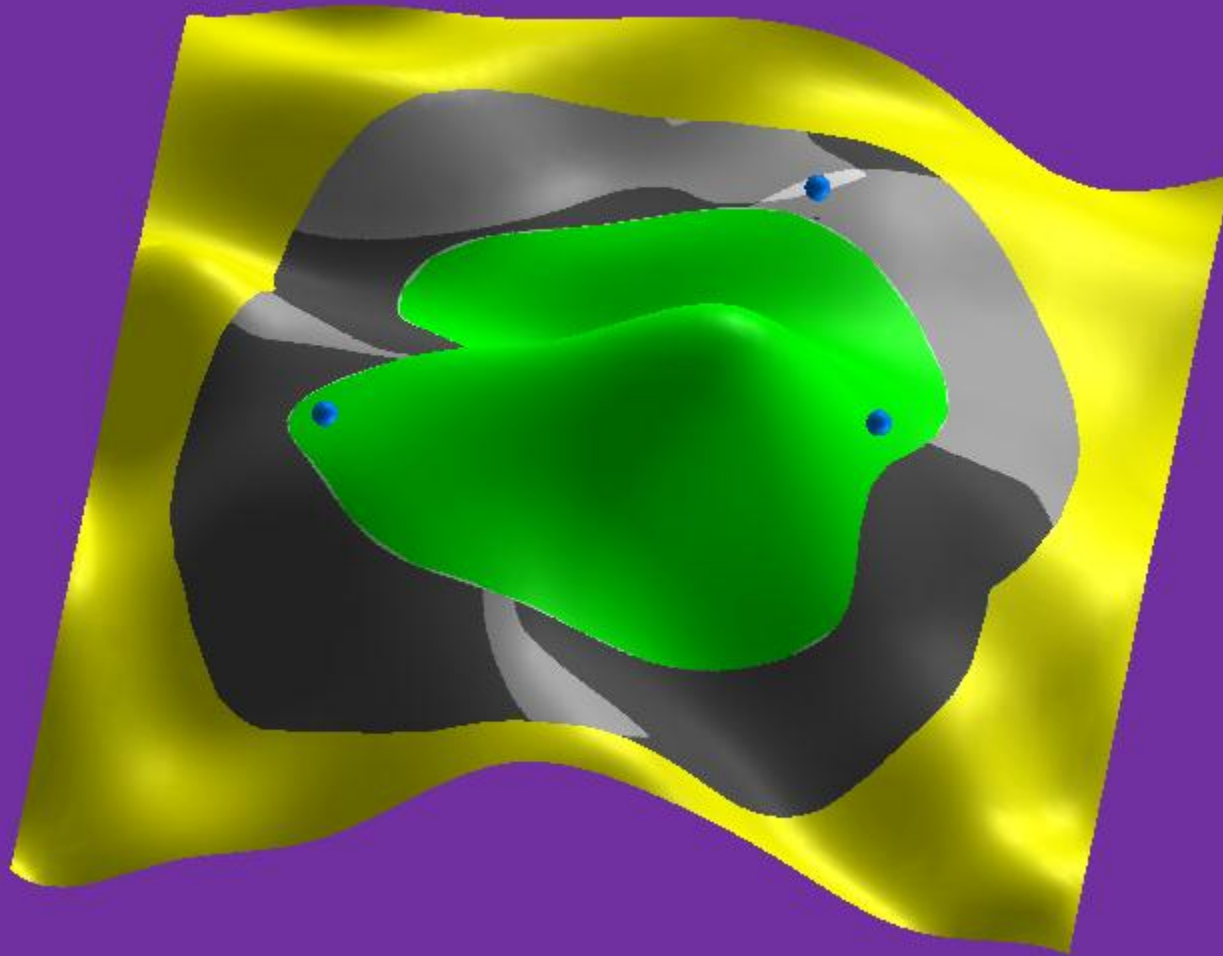
# A military compound



# A military compound - candidates above the perimeter

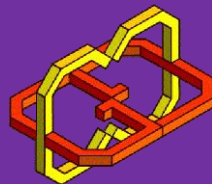
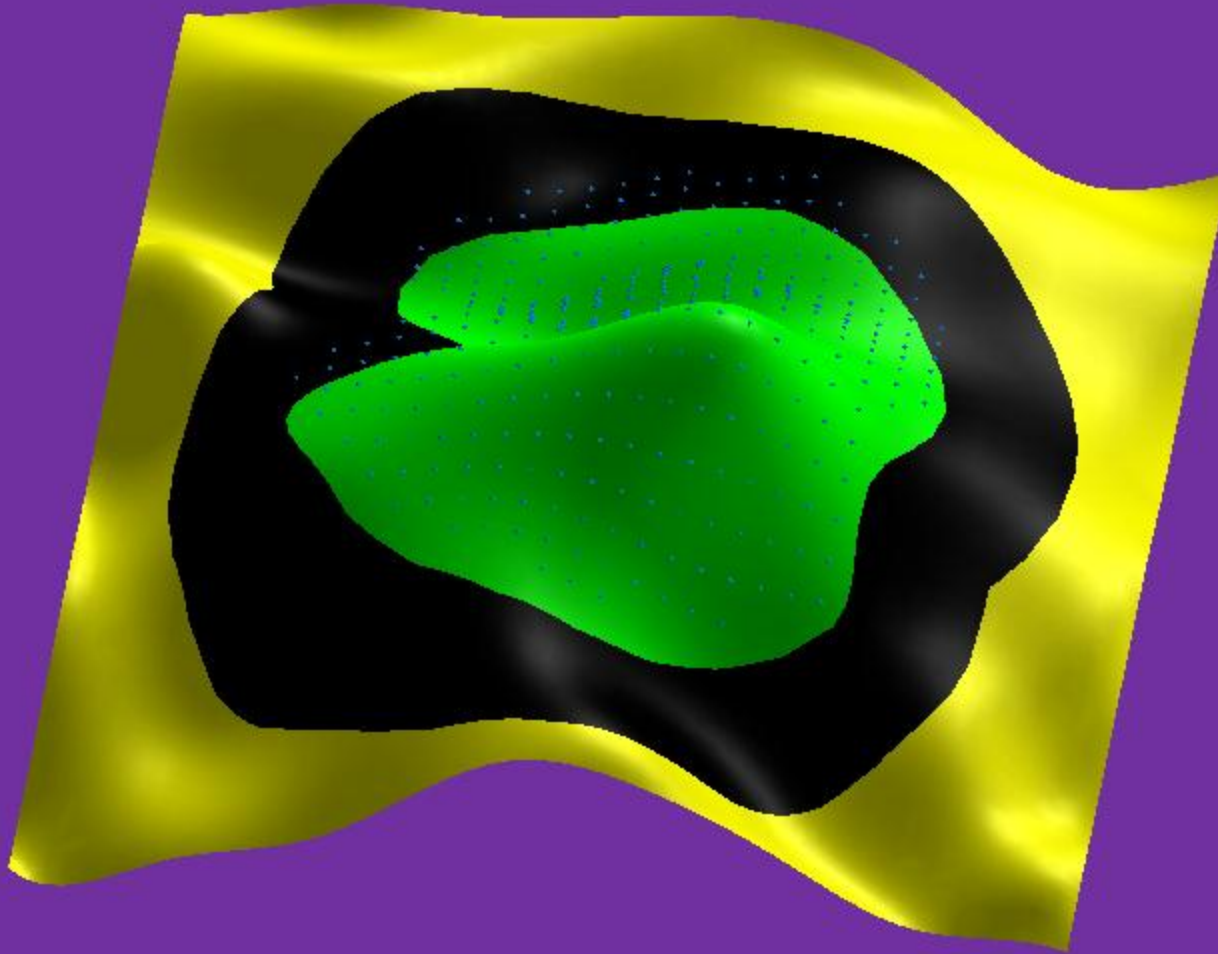


# Candidates above the perimeter – 3 guards solution

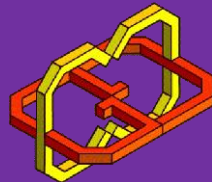
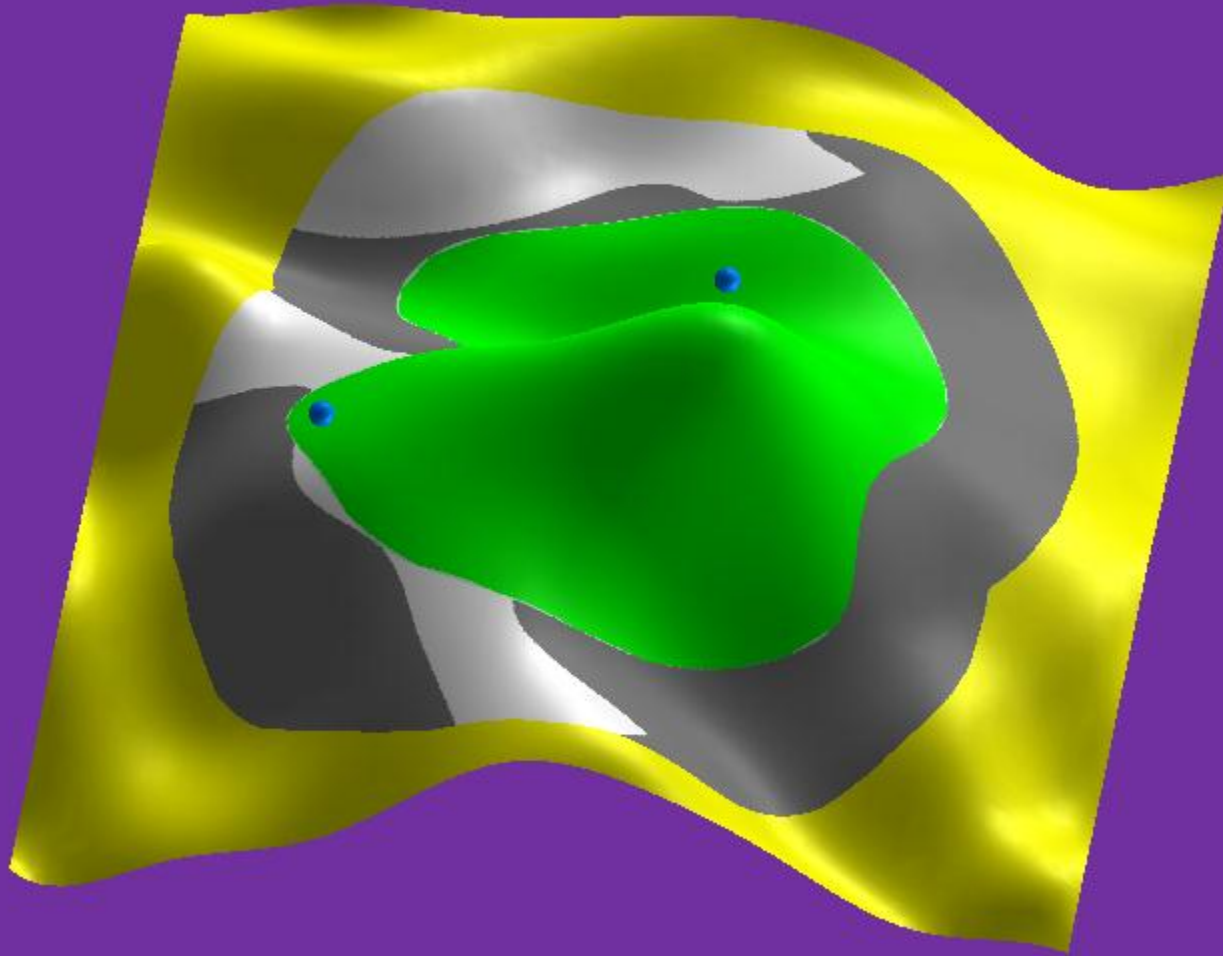




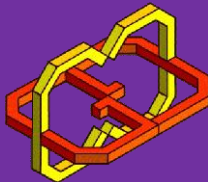
# A military compound - candidates above the compound



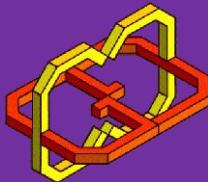
# Candidates above the compound– 2 guards solution



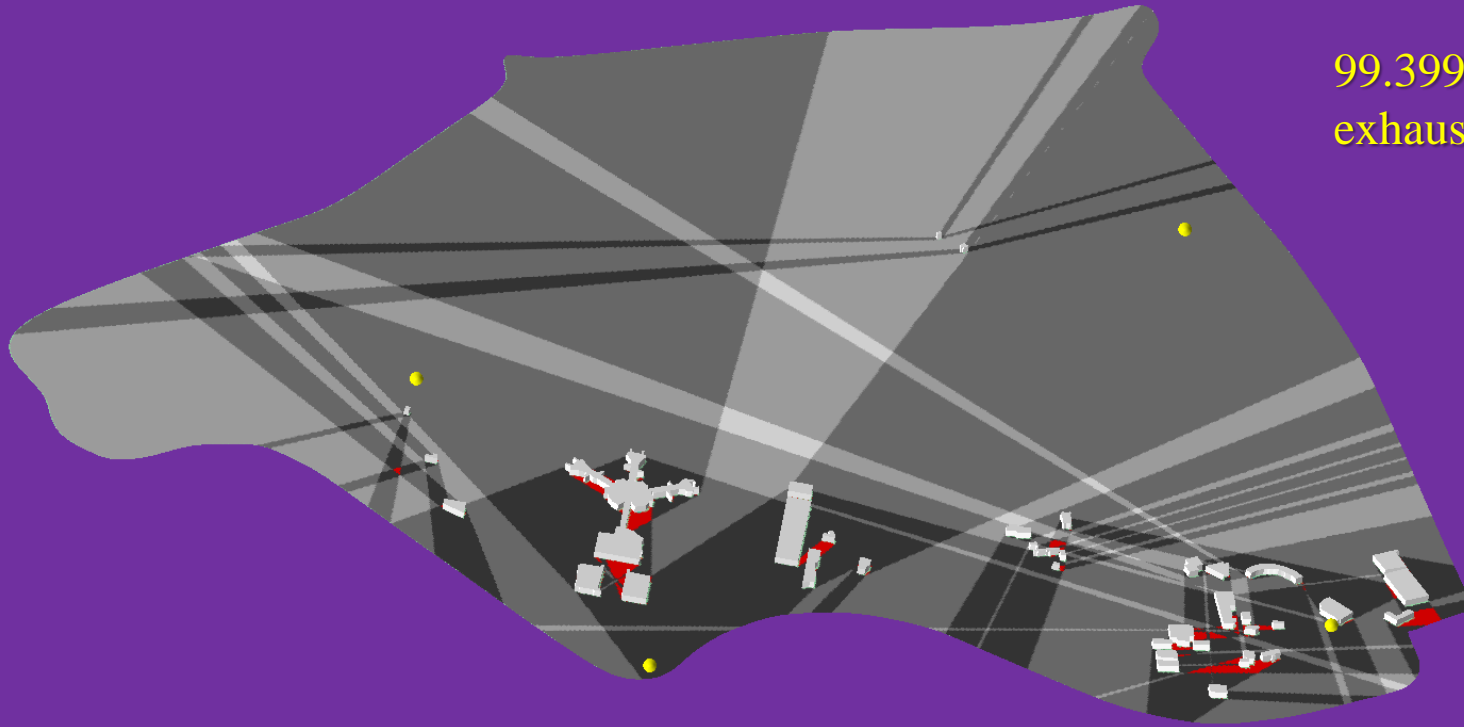
# Ben Gurion airport



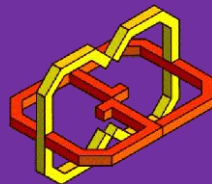
# Ben Gurion airport - candidate cameras



# Ben Gurion airport - exhaustive 4 views solution

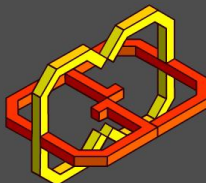


99.399% cover in exhaustive SC.



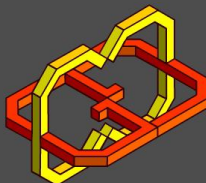
# Conclusions and Future Work I

- ❑ We solve the GC problem in the parametric domain and reduce the analysis into the pixel level.
- ❑ Though we presented the framework in  $R^3$ , nothing prevents the use of this framework in  $R^n$  for arbitrary  $n$ .
- ❑ The reduction to the discrete SC problem allows to optimally solve only discrete GC problems with a few views.
- ❑ We are looking for the solution in the continuous problem.



# Conclusions and Future Work II

- ❑ Use of GPU in proposed framework can benefit the computation times (expect ~two orders of magnitudes).
- ❑ Viewing angle and location distance limitations can be integrated into the creation of the visibility map.
- ❑ Many of the visibility maps are very similar. Can we use this property to reduce set cover calculations?
- ❑ The suggested framework can be used in other GC problems beside mold design and security.





End

