

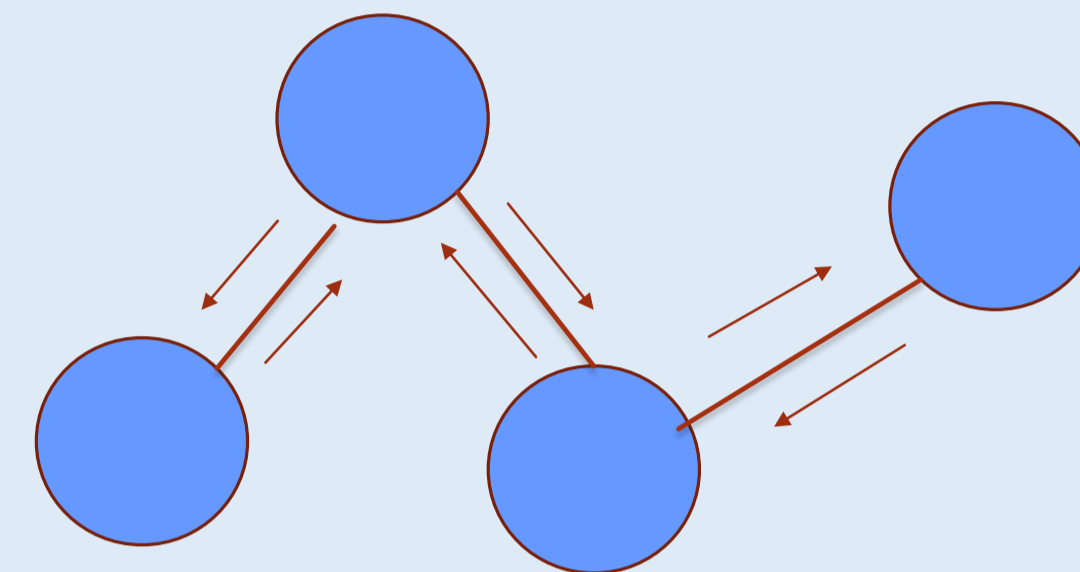
# A fast distributed $2 + \epsilon$ approximation for weighted vertex cover\*

Reuven Bar-Yehuda , Keren Censor-Hillel, Gregory Schwartzman

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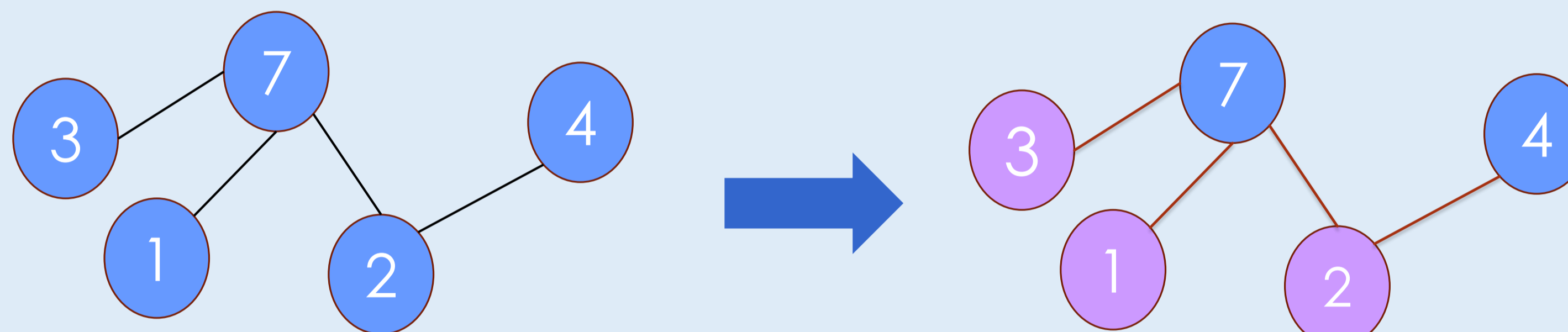
## Model

- A network of computers that send messages in synchronous rounds



## The problem: Vertex cover

- Find a set of nodes of minimal cost that covers all edges

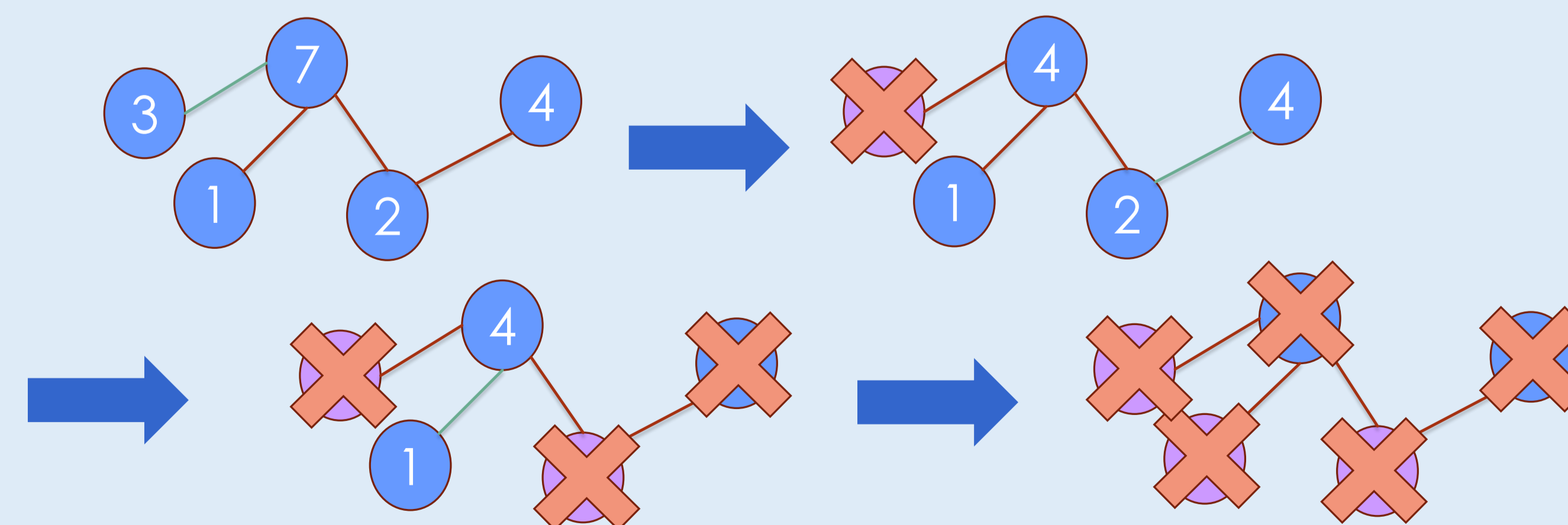


NP-hard, but a 2-approximation exists!

## A sequential 2 approximation

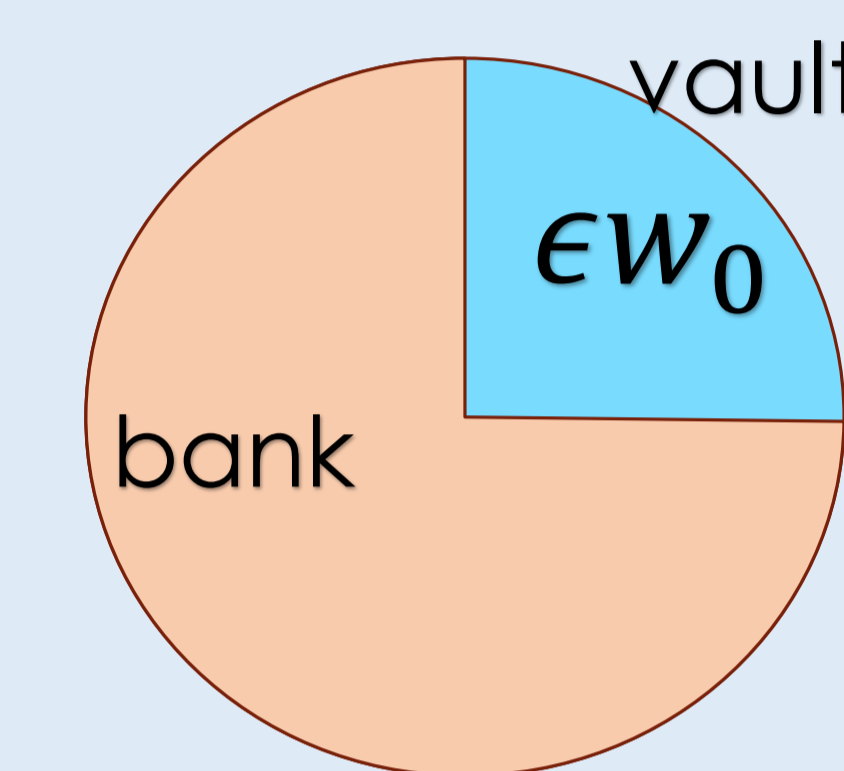
- Choose an edge and remove the minimal weight between its endpoints from both endpoints
- Vertices with weight zero are free, add them!
- Repeat

\*Adding a vertex to the cover when its weight is an  $\epsilon$ -fraction of its original weight results in a  $2+\epsilon$  approximation.



Can be distributed!

## A distributed $(2 + \epsilon)$ approximation

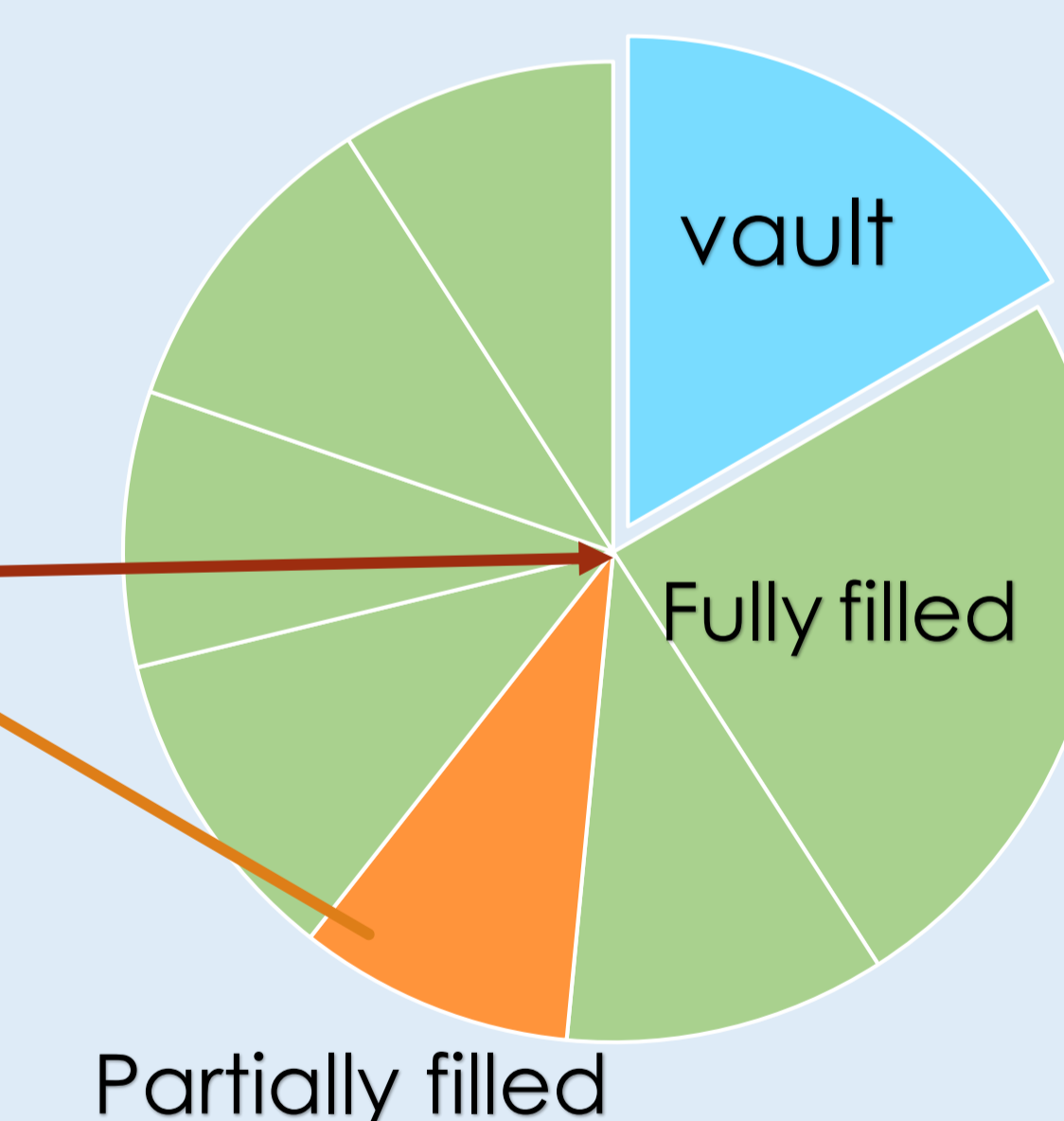
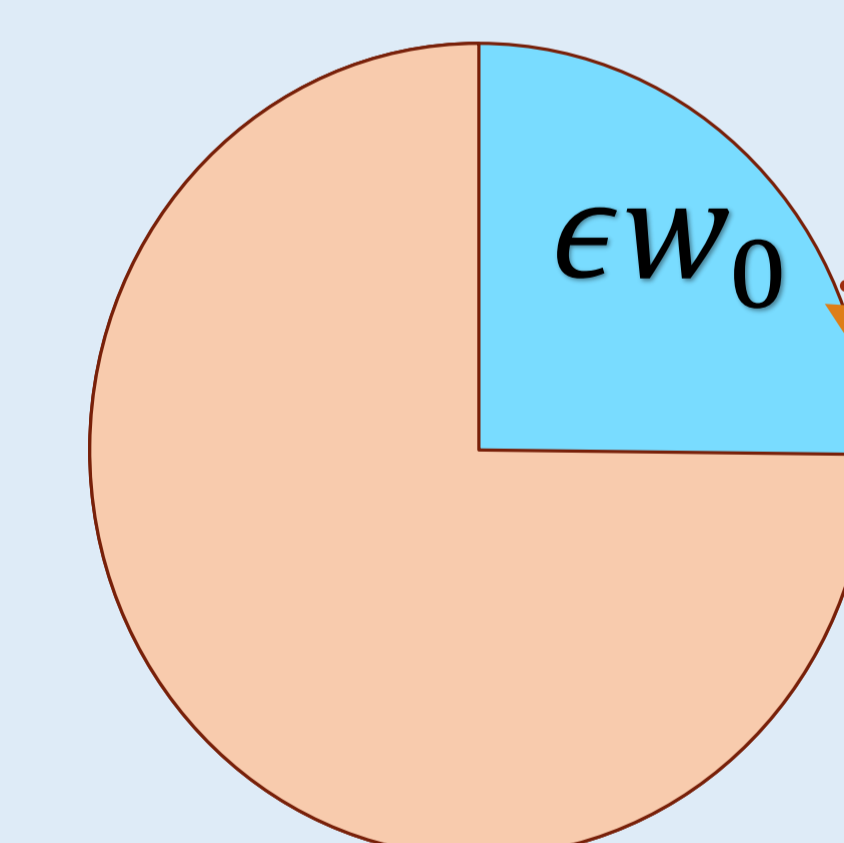


Split the weight of the vertex

Offers are made by vault, bank responds



When the weight goes below  $\epsilon W_0$ , add to cover

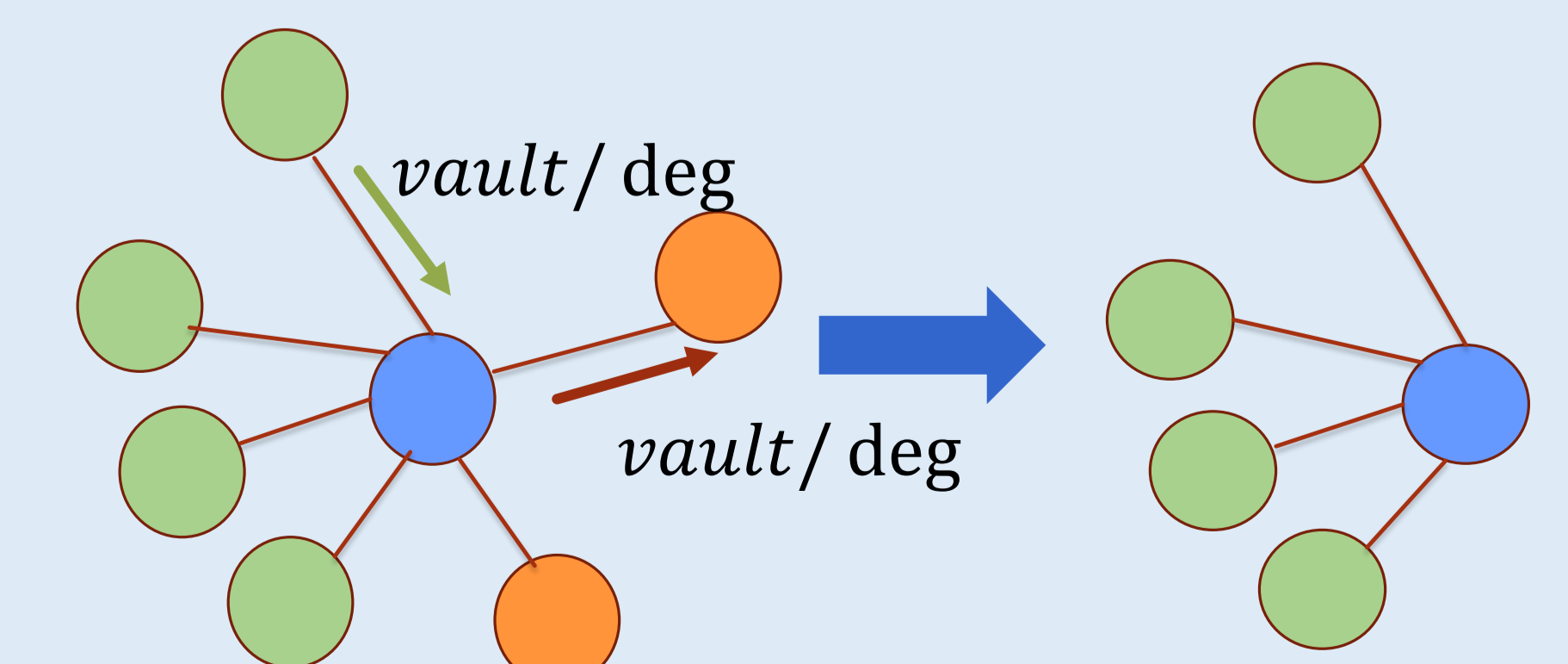


Partially filled

## Analysis

There exists a tradeoff between weight and degree reductions

$O\left(\frac{\log \Delta}{\epsilon \log \log \Delta}\right)$  rounds  
This is tight for constant  $\epsilon$ !



\*Maximal node degree is  $\Delta$